# Forbidden Configurations A shattered history 

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I have had the good fortune of working with a number of coauthors in this area: Farzin Barekat, Laura Dunwoody, Ron Ferguson, Balin Fleming, Zoltan Füredi, Jerry Griggs, Nima Kamoosi, Steven Karp, Peter Keevash, Christina Koch, Linyuan (Lincoln) Lu, Connor Meehan, U.S.R. Murty, Miguel Raggi, Lajos Ronyai, and Attila Sali. A survey paper is available.

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i.e. if $A$ is $m$-rowed then $A$ is the incidence matrix of some family $\mathcal{A}$ of subsets of $[m]=\{1,2, \ldots, m\}$.

$$
\begin{gathered}
A=\left[\begin{array}{lll|l|l}
0 & 0 & 0 & 1 & 1 \\
0 & 1 & 0 & 0 & 1 \\
0 & 0 & 1 & 1 & 1
\end{array}\right] \\
\mathcal{A}=\{\emptyset,\{2\},\{3\},\{1,3\},\{1,2,3\}\}
\end{gathered}
$$

Definition Given a matrix $F$, we say that $A$ has $F$ as a configuration written $F \prec A$ if there is a submatrix of $A$ which is a row and column permutation of $F$.

$$
F=\left[\begin{array}{llll}
0 & 0 & 1 & 1 \\
0 & 1 & 0 & 1
\end{array}\right] \prec\left[\begin{array}{llllll}
1 & 1 & 1 & 1 & 1 & 1 \\
0 & 1 & 0 & 1 & 1 & 0 \\
0 & 0 & 0 & 0 & 1 & 1 \\
0 & 1 & 1 & 0 & 0 & 0
\end{array}\right]=A
$$

## Our Extremal Problem

Definition We define $\|A\|$ to be the number of columns in $A$. Avoid $(m, F)=\{A: A$ is $m$-rowed simple, $F \nprec A\}$

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Example: forb $\left(m,\left[\begin{array}{ll}1 & 0 \\ 0 & 1\end{array}\right]\right)=m+1$.

Definition Let $K_{k}$ denote the $k \times 2^{k}$ simple matrix of all possible columns on $k$ rows.
Theorem (Sauer 72, Perles and Shelah 72, Vapnik and Chervonenkis 71)

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\operatorname{forb}\left(m, K_{k}\right)=\binom{m}{k-1}+\binom{m}{k-2}+\cdots+\binom{m}{0}=\Theta\left(m^{k-1}\right)
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Corollary Let $F$ be a $k \times \ell$ simple matrix. Then forb $(m, F)=O\left(m^{k-1}\right) \quad\left(F \prec K_{k}\right)$

We say a set of rows $S$ is shattered by $A$ if $\left.K_{|S|} \prec A\right|_{S}$. Definition $V C$-dimension $(A)=\max \left\{k: K_{k} \prec A\right\}$

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Definition $V C$-dimension $(A)=\max \left\{k: K_{k} \prec A\right\}$
VC-dimension gets used in Learning Theory and applied probability.

## Let $\operatorname{sh}(A)=\{S \subseteq[m]: A$ shatters $S\}$

e.g.

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& \operatorname{sh}(A)=\{\emptyset,\{1\},\{2\},\{3\},\{4\},\{2,3\},\{2,4\}\} \\
& \text { So }|\operatorname{sh}(A)|=7 \geq 6=\|A\|
\end{aligned}
$$

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Theorem (Pajor 85) $\quad|\operatorname{sh}(A)| \geq\|A\|$.
Proof: Decompose $A$ as follows:

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A=\left[\begin{array}{ccc}
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$|\operatorname{sh}(A)| \geq\left|\operatorname{sh}\left(A_{0}\right)\right|+\left|\operatorname{sh}\left(A_{1}\right)\right|$.
Hence $|\operatorname{sh}(A)| \geq\|A\|$.

Remark If $A$ shatters $S$ then $A$ shatters any subset of $S$.
Theorem (Sauer 72, Perles and Shelah 72, Vapnik and Chervonenkis 71)

$$
\text { forb }\left(m, K_{k}\right)=\binom{m}{k-1}+\binom{m}{k-2}+\cdots+\binom{m}{0}
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Proof: Let $A \in \operatorname{Avoid}\left(m, K_{k}\right)$.
Then $\operatorname{sh}(A)$ can only contain sets of size $k-1$ or smaller.
Then

$$
\binom{m}{k-1}+\binom{m}{k-2}+\cdots+\binom{m}{0} \geq|\operatorname{sh}(A)| \geq\|A\| .
$$

## Critical Substructures

Definition A critical substructure of a configuration $F$ is a minimal configuration $F^{\prime} \prec F$ such that

$$
\text { forb }\left(m, F^{\prime}\right)=\text { forb }(m, F) .
$$

When $F^{\prime} \prec F^{\prime \prime} \prec F$, we deduce that

$$
\text { forb }\left(m, F^{\prime}\right)=\text { forb }\left(m, F^{\prime \prime}\right)=\text { forb }(m, F) \text {. }
$$

Let $1_{k} \mathbf{0}_{\ell}$ denote the $(k+\ell) \times 1$ column of $k$ 's on top of $\ell 0$ 's. Let $K_{k}^{\ell}$ denote the $k \times\binom{ k}{\ell}$ simple matrix of all columns of sum $\ell$.


Miguel Raggi


Steven Karp

## Critical Substructures for $K_{4}$

$$
K_{4}=\left[\begin{array}{llllllllllllllll}
1 & 1 & 1 & 1 & 0 & 1 & 1 & 1 & 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 \\
1 & 1 & 1 & 0 & 1 & 1 & 0 & 0 & 1 & 1 & 0 & 0 & 1 & 0 & 0 & 0 \\
1 & 1 & 0 & 1 & 1 & 0 & 1 & 0 & 1 & 0 & 1 & 0 & 0 & 1 & 0 & 0 \\
1 & 0 & 1 & 1 & 1 & 0 & 0 & 1 & 0 & 1 & 1 & 0 & 0 & 0 & 1 & 0
\end{array}\right]
$$

Critical substructures are $\mathbf{1}_{4}, K_{4}^{3}, K_{4}^{2}, K_{4}^{1}, \mathbf{0}_{4}, 2 \cdot \mathbf{1}_{3}, 2 \cdot \mathbf{0}_{3}$. Note that forb $\left(m, \mathbf{1}_{4}\right)=$ forb $\left(m, K_{4}^{3}\right)=$ forb $\left(m, K_{4}^{2}\right)=$ forb $\left(m, K_{4}^{1}\right)$
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K_{4}=\left[\begin{array}{llllllllllllllll}
1 & 1 & 1 & 1 & 0 & 1 & 1 & 1 & 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 \\
1 & 1 & 1 & 0 & 1 & 1 & 0 & 0 & 1 & 1 & 0 & 0 & 1 & 0 & 0 & 0 \\
1 & 1 & 0 & 1 & 1 & 0 & 1 & 0 & 1 & 0 & 1 & 0 & 0 & 1 & 0 & 0 \\
1 & 0 & 1 & 1 & 1 & 0 & 0 & 1 & 0 & 1 & 1 & 0 & 0 & 0 & 1 & 0
\end{array}\right]
$$

Critical substructures are $\mathbf{1}_{4}, K_{4}^{3}, K_{4}^{2}, K_{4}^{1}, \mathbf{0}_{4}, 2 \cdot \mathbf{1}_{3}, 2 \cdot \mathbf{0}_{3}$. Note that forb $\left(m, \mathbf{1}_{4}\right)=$ forb $\left(m, K_{4}^{3}\right)=$ forb $\left(m, K_{4}^{2}\right)=$ forb $\left(m, K_{4}^{1}\right)$
$=\operatorname{forb}\left(m, \mathbf{0}_{4}\right)=\operatorname{forb}\left(m, 2 \cdot \mathbf{1}_{3}\right)=\operatorname{forb}\left(m, 2 \cdot \mathbf{0}_{3}\right)$.

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1 & 1 & 1 & 0 & 1 & 1 & 0 & 0 & 1 & 1 & 0 & 0 & 1 & 0 & 0 & 0 \\
1 & 1 & 0 & 1 & 1 & 0 & 1 & 0 & 1 & 0 & 1 & 0 & 0 & 1 & 0 & 0 \\
1 & 0 & 1 & 1 & 1 & 0 & 0 & 1 & 0 & 1 & 1 & 0 & 0 & 0 & 1 & 0
\end{array}\right]
$$

Critical substructures are $\mathbf{1}_{4}, K_{4}^{3}, K_{4}^{2}, K_{4}^{1}, \mathbf{0}_{4}, 2 \cdot \mathbf{1}_{3}, 2 \cdot \mathbf{0}_{3}$. Note that forb $\left(m, \mathbf{1}_{4}\right)=$ forb $\left(m, K_{4}^{3}\right)=$ forb $\left(m, K_{4}^{2}\right)=$ forb $\left(m, K_{4}^{1}\right)$
$=$ forb $\left(m, \mathbf{0}_{4}\right)=$ forb $\left(m, 2 \cdot \mathbf{1}_{3}\right)=$ forb $\left(m, 2 \cdot \mathbf{0}_{3}\right)$.
The same is conjectured to be true for $K_{k}$ for $k \geq 5$.

## We can extend $K_{4}$ and yet have the same bound

$\left[K_{4} \mid \mathbf{1}_{2} \mathbf{0}_{2}\right]=$

$$
\left[\begin{array}{llllllllllllllll|l}
1 & 1 & 1 & 1 & 0 & 1 & 1 & 1 & 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 & 1 \\
1 & 1 & 1 & 0 & 1 & 1 & 0 & 0 & 1 & 1 & 0 & 0 & 1 & 0 & 0 & 0 & 1 \\
1 & 1 & 0 & 1 & 1 & 0 & 1 & 0 & 1 & 0 & 1 & 0 & 0 & 1 & 0 & 0 & 0 \\
1 & 0 & 1 & 1 & 1 & 0 & 0 & 1 & 0 & 1 & 1 & 0 & 0 & 0 & 1 & 0 & 0
\end{array}\right]
$$

Theorem (A., Meehan) For $m \geq 5$, we have forb $\left(m,\left[K_{4} \mid \mathbf{1}_{2} \mathbf{0}_{2}\right]\right)=$ forb $\left(m, K_{4}\right)$.

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1 & 1 & 1 & 1 & 0 & 1 & 1 & 1 & 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 & 1 \\
1 & 1 & 1 & 0 & 1 & 1 & 0 & 0 & 1 & 1 & 0 & 0 & 1 & 0 & 0 & 0 & 1 \\
1 & 1 & 0 & 1 & 1 & 0 & 1 & 0 & 1 & 0 & 1 & 0 & 0 & 1 & 0 & 0 & 0 \\
1 & 0 & 1 & 1 & 1 & 0 & 0 & 1 & 0 & 1 & 1 & 0 & 0 & 0 & 1 & 0 & 0
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Theorem (A., Meehan) For $m \geq 5$, we have forb $\left(m,\left[K_{4} \mid \mathbf{1}_{2} \mathbf{0}_{2}\right]\right)=$ forb $\left(m, K_{4}\right)$.
We expect in fact that we could add many copies of the column $\mathbf{1}_{2} \mathbf{0}_{2}$ and obtain the same bound, albeit for larger values of m .


## Connor Meehan

## A Product Construction

The building blocks of our product constructions are $I, I^{c}$ and $T$ :

$$
I_{4}=\left[\begin{array}{llll}
1 & 0 & 0 & 0 \\
0 & 1 & 0 & 0 \\
0 & 0 & 1 & 0 \\
0 & 0 & 0 & 1
\end{array}\right], \quad I_{4}^{c}=\left[\begin{array}{llll}
0 & 1 & 1 & 1 \\
1 & 0 & 1 & 1 \\
1 & 1 & 0 & 1 \\
1 & 1 & 1 & 0
\end{array}\right], \quad T_{4}=\left[\begin{array}{llll}
1 & 1 & 1 & 1 \\
0 & 1 & 1 & 1 \\
0 & 0 & 1 & 1 \\
0 & 0 & 0 & 1
\end{array}\right]
$$

Definition Given an $m_{1} \times n_{1}$ matrix $A$ and a $m_{2} \times n_{2}$ matrix $B$ we define the product $A \times B$ as the $\left(m_{1}+m_{2}\right) \times\left(n_{1} n_{2}\right)$ matrix consisting of all $n_{1} n_{2}$ possible columns formed from placing a column of $A$ on top of a column of $B$. If $A, B$ are simple, then $A \times B$ is simple. (A, Griggs, Sali 97)

$$
\left[\begin{array}{lll}
1 & 0 & 0 \\
0 & 1 & 0 \\
0 & 0 & 1
\end{array}\right] \times\left[\begin{array}{lll}
1 & 1 & 1 \\
0 & 1 & 1 \\
0 & 0 & 1
\end{array}\right]=\left[\begin{array}{lll|lll|lll}
1 & 1 & 1 & 0 & 0 & 0 & 0 & 0 & 0 \\
0 & 0 & 0 & 1 & 1 & 1 & 0 & 0 & 0 \\
0 & 0 & 0 & 0 & 0 & 0 & 1 & 1 & 1 \\
\hline 1 & 1 & 1 & 1 & 1 & 1 & 1 & 1 & 1 \\
0 & 1 & 1 & 0 & 1 & 1 & 0 & 1 & 1 \\
0 & 0 & 1 & 0 & 0 & 1 & 0 & 0 & 1
\end{array}\right]
$$

Given $p$ simple matrices $A_{1}, A_{2}, \ldots, A_{p}$, each of size $m / p \times m / p$, the $p$-fold product $A_{1} \times A_{2} \times \cdots \times A_{p}$ is a simple matrix of size $m \times\left(m^{p} / p^{p}\right)$ i.e. $\Theta\left(m^{p}\right)$ columns.

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0 & 0 & 1
\end{array}\right]=\left[\begin{array}{lll|lll|lll}
1 & 1 & 1 & 0 & 0 & 0 & 0 & 0 & 0 \\
0 & 0 & 0 & 1 & 1 & 1 & 0 & 0 & 0 \\
0 & 0 & 0 & 0 & 0 & 0 & 1 & 1 & 1 \\
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0 & 1 & 1 & 0 & 1 & 1 & 0 & 1 & 1 \\
0 & 0 & 1 & 0 & 0 & 1 & 0 & 0 & 1
\end{array}\right]
$$

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## The Conjecture

Definition Let $x(F)$ denote the largest $p$ such that there is a $p$-fold product which does not contain $F$ as a configuration where the $p$-fold product is $A_{1} \times A_{2} \times \cdots \times A_{p}$ where each $A_{i} \in\left\{I_{m / p}, I_{m / p}^{c}, T_{m / p}\right\}$.

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Conjecture (A, Sali 05) forb $(m, F)$ is $\Theta\left(m^{\times(F)}\right)$.
In other words, we predict our product constructions with the three building blocks $\left\{I, I^{c}, T\right\}$ determine the asymptotically best constructions. The conjecture has been verified in many cases.

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Attila Sali

## An Unavoidable Forbidden Family

Theorem (Balogh and Bollobás 05) Let $k$ be given. Then

$$
\text { forb }\left(m,\left\{I_{k}, I_{k}^{c}, T_{k}\right\}\right) \leq 2^{2^{k}}
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Note that the bound does not depend on $m$ ! Also note that there is no obvious product construction of $I, I^{c}, T$ simultaneously avoiding $I_{k}, I_{k}^{c}, T_{k}$ so this is consistent with the conjecture. It has the spirit of Ramsey Theory.

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Theorem (A., Lu 14) Let $k$ be given. Then there is a constant $c$

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Theorem (A., Lu 14) Let $k$ be given. Then there is a constant $c$

$$
\text { forb }\left(m,\left\{I_{k}, I_{k}^{c}, T_{k}\right\}\right) \leq 2^{c k^{2}}
$$

A construction taking all columns of column sum at most $k-1$ that arise from the $k-1$-fold product $T_{k-1} \times T_{k-1} \times \cdots \times T_{k-1}$ yields forb $\left(m,\left\{I_{k}, I_{k}^{c}, T_{k}\right\}\right) \geq\binom{ 2 k-2}{k-1} \approx 2^{2 k}$.
Probabalistic constructions of Balogh and Bollobás yield forb $\left(m,\left\{I_{k}, I_{k}^{c}, T_{k}\right\}\right) \geq c \cdot 2^{k \log k}$.

The proof uses lots of induction and multicoloured Ramsey numbers: $R\left(k_{1}, k_{2}, \ldots, k_{\ell}\right)$ is the smallest value of $n$ such than any colouring of the edges of $K_{n}$ with $\ell$ colours $1,2, \ldots, \ell$ will have some colour $i$ and a clique of $k_{i}$ vertices with all edges of colour $i$. These numbers are readily bounded by multinomial coefficients:

$$
\begin{gathered}
R\left(k_{1}, k_{2}, \ldots, k_{\ell}\right) \leq\binom{\sum_{i=1}^{\ell} k_{i}}{k_{1} k_{2} k_{3} \cdots k_{\ell}} \\
R\left(k_{1}, k_{2}, \ldots, k_{\ell}\right) \leq 2^{k_{1}+k_{2}+\cdots+k_{\ell}}
\end{gathered}
$$



Linyuan (Lincoln) Lu

As part of our proof we wish to show that we cannot have a $u \times 2 u$ large ( 0,1 )-matrix of the form

| 0 | 1 | $*$ | $*$ | $*$ | $*$ | $*$ | $*$ | $*$ | $*$ | $*$ | $*$ |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| $a$ | $a$ | 0 | 1 | $*$ | $*$ | $*$ | $*$ | $*$ | $*$ | $*$ | $*$ |
| $b$ | $b$ | $c$ | $c$ | 0 | 1 | $*$ | $*$ | $*$ | $*$ | $*$ | $*$ |
| $d$ | $d$ | $e$ | $e$ | $f$ | $f$ | 0 | 1 | $*$ | $*$ | $*$ | $*$ |
| $g$ | $g$ | $h$ | $h$ | $i$ | $i$ | $j$ | $j$ | 0 | 1 | $*$ | $*$ |
| $k$ | $k$ | $l$ | $l$ | $m$ | $m$ | $n$ | $n$ | $o$ | $o$ | 0 | 1 |

One can interpret the entries of the matrix as $1 \times 2$ blocks yielding a $u \times u$ matrix with the blocks below the diagonal either 00 or 11 with blocks on the diagonal 01 and arbitrary ( 0,1 )-blocks above the diagonal.

We consider a colouring of the complete graph $K_{u}$ with edge $i, j$ getting a colour based on the entries in the block $j, i$ and the block $i, j$.

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There are 6 colours to consider

$$
\left[\begin{array}{ll}
0 & 0 \\
0 & 0
\end{array}\right],\left[\begin{array}{ll}
1 & * \\
0 & 0
\end{array}\right],\left[\begin{array}{ll}
* & 1 \\
0 & 0
\end{array}\right],\left[\begin{array}{ll}
1 & 1 \\
1 & 1
\end{array}\right],\left[\begin{array}{ll}
0 & * \\
1 & 1
\end{array}\right],\left[\begin{array}{ll}
* & 0 \\
1 & 1
\end{array}\right]
$$

We are able to show that $u<R(k, k+1, k, k, k+1, k)$ and so we get a singly exponential bound on $u \leq 2^{6 k+3}$. The proof has more to do than this but this is a critical step.

We say that the edge $i, j$ is colour $\left[\begin{array}{ll}1 & * \\ 0 & 0\end{array}\right]$ if we have 00 in entry $(j, i)$ and $1 *$ in entry $(i, j)$ :

$$
\begin{array}{cccc} 
& i & & j \\
i & 01 & & 1 * \\
& & \ddots & \\
j & 00 & & 01
\end{array}
$$

Now consider a clique of size $k+1$ of colour $\left[\begin{array}{ll}1 & * \\ 0 & 0\end{array}\right]$ :

| 0 | 1 | 1 | $*$ | 1 | $*$ | 1 | $*$ | 1 | $*$ |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| 0 | 0 | 0 | 1 | 1 | $*$ | 1 | $*$ | 1 | $*$ |
| 0 | 0 | 0 | 0 | 0 | 1 | 1 | $*$ | 1 | $*$ |
| 0 | 0 | 0 | 0 | 0 | 0 | 0 | 1 | 1 | $*$ |
| 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 1 |

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$$
\begin{array}{cccc} 
& i & & j \\
i & 01 & & 1 * \\
& & \ddots & \\
j & 00 & & 01
\end{array}
$$

Now consider a clique of size $k+1$ of colour $\left[\begin{array}{ll}1 & * \\ 0 & 0\end{array}\right]$ :

$$
\begin{array}{ll|l|l|l|l|l|l|l|ll}
0 & 1 & 1 & * & 1 & * & 1 & * & 1 & * & \\
0 & 0 & 0 & 1 & 1 & * & 1 & * & 1 & * & \\
0 & 0 & 0 & 0 & 0 & 1 & 1 & * & 1 & * & \text { yields } T_{k} \\
0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 & 1 & * & \\
0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1
\end{array}
$$

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& i & & j \\
i & 01 & & * 1 \\
& & \ddots & \\
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\end{array}
$$

Now consider a clique of size $k$ of colour $\left[\begin{array}{ll}* & 1 \\ 0 & 0\end{array}\right]$ :

| 0 | 1 | $*$ | 1 | $*$ | 1 | $*$ | 1 | $*$ | 1 |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| 0 | 0 | 0 | 1 | $*$ | 1 | $*$ | 1 | $*$ | 1 |
| 0 | 0 | 0 | 0 | 0 | 1 | $*$ | 1 | $*$ | 1 |
| 0 | 0 | 0 | 0 | 0 | 0 | 0 | 1 | $*$ | 1 |
| 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 1 |

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$$
\begin{array}{cccc} 
& i & & j \\
i & 01 & & * 1 \\
& & \ddots & \\
j & 00 & & 01
\end{array}
$$

Now consider a clique of size $k$ of colour $\left[\begin{array}{cc}* & 1 \\ 0 & 0\end{array}\right]$ :

$$
\begin{array}{l|l|l|l|l|l|l|l|l|l|}
0 & 1 & * & 1 & * & 1 & * & 1 & * & 1 \\
0 & 0 & 0 & 1 & * & 1 & * & 1 & * & 1 \\
0 & 0 & 0 & 0 & 0 & 1 & * & 1 & * & 1 \\
0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 & * & 1 \\
0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 \\
& &
\end{array} \text { yields } T_{k}
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i & 01 & & 00 \\
& & \ddots & \\
j & 00 & & 01
\end{array}
$$

Now consider a clique of size $k$ of colour $\left[\begin{array}{ll}0 & 0 \\ 0 & 0\end{array}\right]$ :

| 0 | 1 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| 0 | 0 | 0 | 1 | 0 | 0 | 0 | 0 | 0 | 0 |
| 0 | 0 | 0 | 0 | 0 | 1 | 0 | 0 | 0 | 0 |
| 0 | 0 | 0 | 0 | 0 | 0 | 0 | 1 | 0 | 0 |
| 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 1 |

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& i & & j \\
i & 01 & & 00 \\
& & \ddots & \\
j & 00 & & 01
\end{array}
$$

Now consider a clique of size $k$ of colour $\left[\begin{array}{ll}0 & 0 \\ 0 & 0\end{array}\right]$ :

$$
\begin{array}{l|l|l|l|l|l|l|l|l|l|}
0 & 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\
0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 & 0 & 0 \\
0 & 0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 \\
0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 & 0 & 0 \\
0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 \\
\cline { 2 - 5 } & &
\end{array}
$$

## Design Theory

Let $s \cdot F$ denote $\overbrace{[F|F| \cdots \mid F]}^{s}$.
Theorem (A, Füredi 86) Let $s \geq 2$ be given. Then forb $\left(m, s \cdot K_{k}\right)$ is $\Theta\left(m^{k}\right)$.

Corollary (Füredi 83) Let $F$ be a $k \times \ell(0,1)$-matrix. Then forb $(m, F)$ is $\Theta\left(m^{k}\right)$.

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Corollary (Füredi 83) Let $F$ be a $k \times \ell(0,1)$-matrix. Then forb $(m, F)$ is $\Theta\left(m^{k}\right)$.

Theorem (A, Sali 14) Let $\alpha$ be given. forb $\left(m, m^{\alpha} \cdot K_{k}\right)$ is $\Theta\left(m^{k+\alpha}\right)$

Note that we are having $s$ grow with $m$. Our forbidden configuration is not fixed but depends on $m$.

Theorem forb $\left(m, m \cdot\left[\begin{array}{l}1 \\ 1\end{array}\right]\right)=\binom{m}{0}+\binom{m}{1}+\binom{m}{2}+\binom{m}{3}$.
Proof: We note that $\left[K_{m}^{0} K_{m}^{1} K_{m}^{2} K_{m}^{3}\right] \in \operatorname{Avoid}\left(m, m \cdot\left[\begin{array}{l}1 \\ 1\end{array}\right]\right)$.
Thus forb $\left(m, m \cdot\left[\begin{array}{l}1 \\ 1\end{array}\right]\right) \geq\binom{ m}{0}+\binom{m}{1}+\binom{m}{2}+\binom{m}{3}$. (note that each pair of rows of has $(m-1) \cdot\left[\begin{array}{l}1 \\ 1\end{array}\right]$ )
We can argue, using the pigeonhole argument,

$$
\text { forb }\left(m, m \cdot\left[\begin{array}{l}
1 \\
1
\end{array}\right]\right) \leq\binom{ m}{0}+\binom{m}{1}+\binom{m}{2}+\frac{m-2}{3}\binom{m}{2}
$$

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1
\end{array}\right]\right) \leq\binom{ m}{0}+\binom{m}{1}+\binom{m}{2}+\binom{m}{3}
$$

Thus forb $\left(m, m \cdot\left[\begin{array}{l}1 \\ 1\end{array}\right]\right)$ is $\Theta\left(m^{3}\right)$.
Can we deduce the growth of forb $\left(m, m^{\alpha} \cdot\left[\begin{array}{l}1 \\ 1\end{array}\right]\right)$ ?

Theorem
forb $\left(m,\left(m+\binom{m-2}{2}\right) \cdot\left[\begin{array}{l}1 \\ 1\end{array}\right]\right)=\binom{m}{0}+\binom{m}{1}+\binom{m}{2}+\binom{m}{3}+\binom{m}{4}$.
$\operatorname{Note}\left[K_{m}^{0} K_{m}^{1} K_{m}^{2} K_{m}^{3} K_{m}^{4}\right] \in \operatorname{Avoid}\left(m,\left(m+\binom{m-2}{2}\right) \cdot\left[\begin{array}{l}1 \\ 1\end{array}\right]\right)$.

Definition Given integers $k \leq m$, let $\binom{[m]}{k}$ denote all $k$ - subsets of [ $m$ ].
Definition Given parameters $t, m, k, \lambda$, a $t-(m, k, \lambda)$ design $\mathcal{D}$ is a multiset of subsets in $\binom{[m]}{k}$ such that for each $S \in\binom{[m]}{t}$ there are exactly $\lambda$ blocks $B \in \mathcal{D}$ containing $S$.
Definition $\operatorname{t} t(m, k, \lambda)$ design $\mathcal{D}$ is simple if $\mathcal{D}$ is a set (i.e. no repeated blocks).

If we have a $t-(m, k, \lambda)$ simple design $\mathcal{D}$, then we can form a matrix $M$ as the element-block incidence matrix associated with $\mathcal{D}$ and we deduce that

$$
\|M\|=\left(\lambda\binom{m}{k} /\binom{k}{t}\right) \text { and } M \in \operatorname{Avoid}\left(m,(\lambda+1) \cdot \mathbf{1}_{t}\right)
$$

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\left.\lambda\binom{m}{k} /\binom{k}{t}\right) \text { and } M \in \operatorname{Avoid}\left(m,(\lambda+1) \cdot \mathbf{1}_{t}\right), ~(1)
\end{array}\right.
$$

$M$ has all columns of sum $k$. We can extend $M$ :

$$
A=\left[K_{m}^{0} K_{m}^{1} K_{m}^{2} \cdots K_{m}^{k-1} M\right]
$$

If we let $\mu=\binom{m-t}{0}+\binom{m-t}{1}+\cdots\binom{m-t}{k-1-t}+\lambda+1$, then $A \in \operatorname{Avoid}\left(m, \mu \cdot \mathbf{1}_{t}\right)$

We can deduce that
$\operatorname{forb}\left(m, \mu \cdot \mathbf{1}_{t}\right)=\binom{m}{0}+\binom{m}{1}+\binom{m}{2}+\cdots+\binom{m}{k-1}+\lambda\binom{m}{k} /\binom{k}{t}$

## Breakthrough of Keevash on the Existence of Designs

Theorem (Keevash 14) Let $1 / m \ll \theta \ll 1 / k \leq 1 /(t+1)$ and $\theta \ll 1$. Suppose that $\binom{k-i}{t-i}$ divides $\binom{m-i}{t-i}$ for $0 \leq i \leq r-1$. Then there exists a $t-(m, k, \lambda)$ simple design for $\lambda \leq \theta m^{k-t}$.

Corollary (Weak Packing Idea) forb $\left(m, m^{\alpha} \cdot \mathbf{1}_{k}\right)$ is $\Theta\left(m^{k+\alpha}\right)$.
For our purposes we don't care about equality but merely asymptotics. We use the Keevash result to establish lower bounds. His result is the first to establish this (of course his results do much more!).

Using the shifting idea, we have

$$
\operatorname{forb}\left(m, s \cdot K_{k}\right)=\operatorname{forb}\left(m, s \cdot \mathbf{1}_{k}\right)
$$

And this establishes the result:
Theorem (A., Sali 14) Let $\alpha$ be given. forb $\left(m, m^{\alpha} \cdot K_{k}\right)$ is $\Theta\left(m^{k+\alpha}\right)$

## Main Upper Bound Proof

Lemma Let $F$ be a simple matrix and let $s>1$ be given. forb $(m, s \cdot F) \leq \sum_{i=1}^{m-1}(s-1) \cdot$ forb $(m-i, F)$

Proof: We use the induction idea of $A$. and Lu 13. The idea is to temporarily allow the matrices to be non-simple in a restricted way.

$$
\text { Let } F=\left[\begin{array}{ll}
1 & 1 \\
1 & 0 \\
0 & 1 \\
0 & 0
\end{array}\right]
$$

We have forb $(m, F)=4 m$, i.e. forb $(m, F)$ is $O(m)$.

Theorem Let $\alpha>0$ be given. Using the Weak Packing idea, forb $\left(m, m^{\alpha} \cdot F\right)$ is $\Theta\left(m^{2+\alpha}\right)$.
Proof:
forb $\left(m, m^{\alpha} \cdot F\right) \leq \sum_{i=1}^{m-1} m^{\alpha} \cdot$ forb $(m-i, F)=m^{\alpha} \sum_{i=1}^{m-1} 4(m-i)$.
Now $\left[\begin{array}{l}1 \\ 1\end{array}\right] \prec F$ and so $m^{\alpha} \cdot\left[\begin{array}{l}1 \\ 1\end{array}\right] \prec m^{\alpha} \cdot F$ from which we have
forb $\left(m, m^{\alpha} \cdot F\right) \geq \operatorname{forb}\left(m, m^{\alpha} \cdot\left[\begin{array}{l}1 \\ 1\end{array}\right]\right)$.

## An Open Problem

$$
\text { Let } F=\left[\begin{array}{ll}
1 & 0 \\
1 & 0 \\
0 & 1 \\
0 & 1
\end{array}\right]
$$

Theorem (Frankl, Füredi, Pach 87) forb $(m, F)=\binom{m}{2}+2 m-1$ i.e. forb $(m, F)$ is $O\left(m^{2}\right)$.

Theorem (A. and Lu 13) Let $s$ be given. Then forb $(m, s \cdot F)$ is $\Theta\left(m^{2}\right)$.
Conjecture forb $\left(m, m^{\alpha} \cdot F\right)$ is $\Theta\left(m^{2+\alpha}\right)$.
We can only prove that forb $\left(m, m^{\alpha} \cdot F\right)$ is $O\left(m^{3+\alpha}\right)$.

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