

DIGRAPHS, DETAILS, AND ACYCLIC ORIENTATIONS

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Math 442-201 2019WT2

31 March 2020

DIRECTED GRAPH

DEFINITION

A **directed graph** (or digraph) D has vertices $V(D)$ and directed **arcs** $A(D)$ connecting the vertices. It is **simple** if

- finite vertices
- no arc joining vertex to itself
- not multiple arcs **going in same direction**.

Note: Every arc **must** have a direction.

Example

$D =$ is a simple digraph.

UNDERLYING GRAPH

DEFINITION

Let D be a digraph. Then the **underlying graph** of D , $U(D)$, is the graph obtained by removing all the arc directions.

Note: $U(D)$ simple $\Rightarrow D$ simple, but D simple $\nRightarrow U(D)$ simple.

Example

$D =$ is a simple digraph with $U(D) =$

ISOMORPHIC

DEFINITION

Let D_1 and D_2 be two digraphs. They are **isomorphic** if and only if $U(D_1) \cong U(D_2)$ **and** the isomorphism preserves the arc orderings.

Example

DEGREES

DEFINITION

Let D be a digraph. The number of arc ends meeting at a vertex v is called the **degree**, $\deg(v)$. Arcs coming in is the **indegree**, $\text{indeg}(v)$. Arcs going out is the **outdegree**, $\text{outdeg}(v)$.

Note:

- $\deg(v) = \text{indeg}(v) + \text{outdeg}(v)$.
- $\text{indeg}(v) = 0$ is a **source**, $\text{outdeg}(v) = 0$ is a **sink**.

Example

$D =$ has $\text{indeg}(v) =$, $\text{outdeg}(v) =$, $\deg(v) =$

ADJACENCY MATRIX

DEFINITION

Let D be a digraph with n vertices v_1, v_2, \dots, v_n . If two vertices are joined by an arc they are **adjacent**. The **adjacency matrix** $A = (a_{ij})$ of D has

$$a_{ij} = \text{no. arcs from } v_i \text{ to } v_j.$$

Note: Vertex v_i is a

- sink if **row** i is all 0,
- source if **column** i is all 0.

Example

$$D = \quad \text{has } A = \begin{pmatrix} & & \\ & & \\ & & \end{pmatrix}$$

(STRONGLY) CONNECTED

Many concepts are similar to graphs, e.g. a **directed path/cycle** is a path/cycle whose directions you follow.

DEFINITION

Let D be a digraph. It is **connected** if $U(D)$ is connected. It is **strongly connected** if it is connected **and** a directed path between every two vertices in **both** directions.

Example

$D =$

$\tilde{D} =$

is connected but **not** strongly.

is both connected **and** strongly.

ORIENTABLE

DEFINITION

Let G be a graph. If we can turn edges into arcs so digraph is strongly connected, then G is **orientable**.

Note: If G is **Eulerian** then G is orientable - just mark the direction of travel on the Eulerian trail.

Example

$G =$

is orientable.

$H =$

is not orientable.

ORIENTATIONS

DEFINITION

Let G be a graph. If we turn edges into arcs, then the digraph D_G is an **orientation** of G .

If D_G has no directed cycles, then it is an **acyclic orientation**. The number of acyclic orientations is $a(G)$.

Example Orientations of K_3 are below, so $a(G) =$

ACYCLIC ORIENTATIONS



THEOREM (STANLEY, 1973)

Let G be a simple graph with n vertices. Then

MIT

Proof idea About 30 mins, using a deletion/contraction for $a(G)$, and then an induction.

Example Let $G = K_3$.

IN SUMMARY

- Digraphs and underlying graphs
- Details: isomorphic, adjacency matrix, directed paths and cycles
- Connected and strongly connected
- Orientable and Eulerian graphs
- Enumerating acyclic orientations

Thanks! See you next time!