

Design Theory and Some Non-simple Forbidden Configurations

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Abstract

Let $\mathbf{1}_k\mathbf{0}_l$ denote the $(k+l) \times 1$ column of k 1's above l 0's. Let $q \cdot (\mathbf{1}_k\mathbf{0}_l)$ denote the $(k+l) \times q$ matrix with q copies of the column $\mathbf{1}_k\mathbf{0}_l$. A 2-design $S_\lambda(2, 3, v)$ can be defined as a $v \times \frac{\lambda}{3} \binom{v}{2}$ (0,1)-matrix with all column sums equal 3 and with no submatrix $(\lambda+1) \cdot (\mathbf{1}_2\mathbf{0}_0)$. Consider an $m \times n$ matrix A with all column sums in $\{3, 4, \dots, m-1\}$. Assume m is sufficiently large (with respect to λ) and assume that A has no submatrix which is a row permutation of $(\lambda+1) \cdot (\mathbf{1}_2\mathbf{0}_1)$. Then the number of columns in A is at most $\frac{\lambda}{3} \binom{m}{3}$ with equality for A being the columns of column sum 3 corresponding to the triples of a 2-design $S_\lambda(2, 3, m)$.

Define a matrix to be *simple* if it is a (0,1)-matrix with no repeated columns. Given two matrices A, F , we define A to have F as a *configuration* if and only if some submatrix of A is a row and column permutation of F . Given m , let $\text{forb}(m, q \cdot (\mathbf{1}_k\mathbf{0}_l))$ denote the maximum number of possible columns in a simple m -rowed matrix which has no configuration $q \cdot (\mathbf{1}_k\mathbf{0}_l)$. For m sufficiently large with respect to q , we compute exact values for $\text{forb}(m, q \cdot (\mathbf{1}_1\mathbf{0}_1))$, $\text{forb}(m, q \cdot (\mathbf{1}_2\mathbf{0}_1))$, $\text{forb}(m, q \cdot (\mathbf{1}_2\mathbf{0}_2))$. In the latter two cases, we use a construction of Dehon (1983) of *simple* triple systems $S_\lambda(2, 3, v)$ for $\lambda > 1$. Moreover for $l = 1, 2$, $m \times \text{forb}(m, q \cdot (\mathbf{1}_2\mathbf{0}_l))$ simple matrices with no configuration $q \cdot (\mathbf{1}_2\mathbf{0}_l)$ must arise from simple 2-designs $S_\lambda(2, 3, m)$ of appropriate λ .

The proofs derive a basic upper bound by a pigeonhole argument and then use careful counting and Turán's bound, for large m , to reduce the bound. For small m , the larger pigeonhole bounds are sometimes the exact bound. There are intermediate values of m for which we do not know the exact bound.

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1 Introduction

Some combinatorial objects can be defined by forbidden substructures. It is also true that most combinatorial objects can be encoded by a (0,1)-matrix. In this paper we consider submatrices of (0,1)-matrices as the substructures of interest.

Let $\mathbf{1}_k\mathbf{0}_l$ denote the $(k+l) \times 1$ column consisting of k 1's atop l 0's. For any positive integer q , let $q \cdot (\mathbf{1}_k\mathbf{0}_l)$ denote the $q \times (k+l)$ matrix of q copies of $\mathbf{1}_k\mathbf{0}_l$. A 2-design $S_\lambda(2, 3, v)$ consists of $\frac{\lambda}{3}\binom{v}{2}$ triples from $\{1, 2, \dots, v\}$ such that for each pair $i, j \in \{1, 2, \dots, v\}$, there are exactly λ triples containing i, j . If we encode the triple system as a v -rowed (0,1)-matrix A such that the columns are the incidence vectors of the triples, then A has no submatrix $(\lambda+1) \cdot (\mathbf{1}_2\mathbf{0}_0)$. In fact, if A is a $v \times n$ (0,1)-matrix with column sums 3 and A has no submatrix $(\lambda+1) \cdot (\mathbf{1}_2\mathbf{0}_0)$ then $n \leq \frac{\lambda}{3}\binom{v}{2}$ with equality if and only if the columns of A correspond to the triples of a 2-design $S_\lambda(2, 3, v)$. This can be shown by a pigeonhole counting argument.

The problem of forbidding a submatrix is usually extended to forbidding any row and column permutation of the submatrix. Let A, F be (0,1)-matrices. We say that A has F as a *configuration* if there is a submatrix of A which is a row and column permutation of F . We extend the forbidden submatrix $(\lambda+1) \cdot (\mathbf{1}_2\mathbf{0}_0)$ and obtain the following two design theory results.

Theorem 1.1 *Let λ and v be given integers. There exists an M so that for $v > M$, if A is an $v \times n$ (0,1)-matrix with column sums in $\{3, 4, \dots, v-1\}$ and A has no configuration $(\lambda+1) \cdot (\mathbf{1}_2\mathbf{0}_1)$ then*

$$n \leq \frac{\lambda}{3} \binom{v}{2} \tag{1}$$

and we have equality if and only if the columns of A correspond to the triples of a 2-design $S_\lambda(2, 3, v)$. ■

When we extend the forbidden configuration to $(\lambda+1) \cdot (\mathbf{1}_2\mathbf{0}_2)$ the case of equality becomes more difficult.

Theorem 1.2 *Let λ and v be given integers. There exists an M so that for $v > M$, if A is an $v \times n$ (0,1)-matrix with column sums in $\{3, 4, \dots, v-3\}$ and A has no configuration $(\lambda+1) \cdot (\mathbf{1}_2\mathbf{0}_2)$ then*

$$n \leq \frac{\lambda}{3} \binom{v}{2} \tag{2}$$

and we have equality if and only if there are positive integers a, b satisfying $a+b = \lambda$ and there are $\frac{a}{3}\binom{v}{2}$ columns of A of column sum 3 corresponding to the triples of a 2-design $S_a(2, 3, v)$ and there are $\frac{b}{3}\binom{v}{2}$ columns of A of column sum $v-3$ of $v-3$ -sets whose complements (in $\{1, 2, \dots, v\}$) corresponding to the triples of a 2-design $S_b(2, 3, v)$. ■

Our first motivation for studying these problems came from extremal set theory. An $m \times n$ (0,1)-matrix A can be thought of a multiset of n subsets of $\{1, 2, \dots, m\}$. Let $[m] = \{1, 2, \dots, m\}$. For an $m \times 1$ (0,1)-column α , we define

$$S(\alpha) = \{i \in [m] : \alpha \text{ has } 1 \text{ in row } i\}. \quad (3)$$

From this we define the natural multiset system \mathcal{A} associated with the matrix A :

$$\mathcal{A} = \{S(\alpha_i) : \alpha_i \text{ is column } i \text{ of } A\}. \quad (4)$$

Similarly, if we are given a multiset system \mathcal{A} , we can form a matrix A , as long as we don't care about column order. We define a *simple* matrix A as a (0,1)-matrix with no repeated columns. In this case \mathcal{A} yields a *set* system and it is in this setting that extremal set theory problems can be stated.

We define $\text{forb}(m, F)$ as the smallest value (depending on m and F) so that if A is a simple $m \times n$ matrix and A has no configuration F then $n \leq \text{forb}(m, F)$. Alternatively $\text{forb}(m, F)$ is the smallest value so that if A is an $m \times (\text{forb}(m, F) + 1)$ simple matrix then A must have a configuration F . A sampling of exact results for $\text{forb}(m, F)$ are in [1], [2].

Let K_k denote the $k \times 2^k$ simple matrix of all possible (0,1)-columns on k rows and let K_k^s denote the $k \times \binom{k}{s}$ simple matrix of all possible columns of column sum s . Many results have been obtained about $\text{forb}(m, F)$. Exact results have been rare for non-simple configurations F . We consider $F = q \cdot (\mathbf{1}_k \mathbf{0}_l)$ for $(k, l) = (1, 1), (2, 1), (2, 2)$. In [1] we showed that

$$\left\lfloor \frac{q+1}{2} m \right\rfloor + 2 \leq \text{forb}(m, q \cdot (\mathbf{1}_1 \mathbf{0}_1)) \leq \left\lfloor \frac{q+1}{2} m + \frac{(q-3)m}{2(m-2)} \right\rfloor + 2$$

where the upper bound obtained by a pigeonhole argument is achieved for $m = q - 1$ by taking $A = [K_m^0 K_m^1 K_m^2 K_m^{m-1} K_m^m]$. For m with $m \geq \max\{3q + 2, 8q - 19\}$, we are able to show that the lower bound is correct and slice $\frac{(q-3)m}{2(m-2)} \approx \frac{q-3}{2}$ off the pigeonhole bound. It is likely that our bound is valid for smaller $m > q - 1$. The case $q = 4$, is Lemma 3.1 in [2] and took a page to establish.

Theorem 1.3 *Let $q \geq 3$ be given. Then for $m \geq \max\{3q + 2, 8q - 19\}$,*

$$\text{forb}(m, q \cdot (\mathbf{1}_1 \mathbf{0}_1)) = \overbrace{\left[\begin{array}{cccc} 1 & 1 & \cdots & 1 \\ 0 & 0 & \cdots & 0 \end{array} \right]}^q = \left\lfloor \frac{q+1}{2} m \right\rfloor + 2. \quad \blacksquare \quad (5)$$

For m even or $q - 3$ even, let G be a (simple) graph on m vertices for which all the degrees are $q - 3$ and for $m, q - 3$ odd let G be a graph for which $m - 1$ vertices have degree $q - 3$ and one vertex has degree $q - 4$. Such graphs are easy to construct. Let H be the vertex-edge incidence matrix associated with G , namely for each edge $e = (i, j)$

of G , we add a column to H with 1's in rows i, j and 0's in other rows. Thus H is a simple m -rowed matrix with $\lfloor \frac{(q-3)m}{2} \rfloor$ columns each of column sum 2. The simple matrix $A = [K_m^0 K_m^1 H K_m^{m-1} K_m^m]$ has $\lfloor \frac{(q+1)m}{2} \rfloor + 2$ columns and no configuration $q \cdot (\mathbf{1}_1 \mathbf{0}_1)$ which establishes $\text{forb}(m, q \cdot (\mathbf{1}_1 \mathbf{0}_1)) \geq \lfloor \frac{(q+1)m}{2} \rfloor + 2$. We establish the upper bound in Section 2.

We are able to solve two more cases but need certain designs to achieve exact bounds. A 2-design $S_\lambda(2, 3, v)$ (or *triple system*) is defined to be *simple* if no triple is repeated. The associated $v \times \frac{\lambda}{3} \binom{v}{2}$ matrix is a simple matrix. We need the following result.

Theorem 1.4 *Dehon[3] Let v, λ be given. Then a simple 2-design $S_\lambda(2, 3, v)$ exists if and only if $v(v-1) \equiv 0 \pmod{6}$, $v-1 \equiv 0 \pmod{2}$ and $v \geq \lambda + 2$. ■*

These designs are used in the constructions for the following two theorems in the following way. We form a simple $v \times \frac{\lambda}{3} \binom{v}{2}$ matrix $T_{v,\lambda}$ whose columns correspond to the blocks of $S_\lambda(2, 3, v)$ so that if B is a block then the corresponding column has a 1 in row i if and only if $i \in B$. Note that $T_{v,\lambda}$ has no submatrix $(\lambda+1) \cdot (\mathbf{1}_2 \mathbf{0}_0)$. Pigeonhole arguments will show that $\text{forb}(m, q \cdot (\mathbf{1}_2 \mathbf{0}_0)) \leq \binom{m}{0} + \binom{m}{1} + \frac{q+1}{3} \binom{m}{2}$ with equality, by Dehon's Theorem 1.4, for $m \geq q$ and $m \equiv 1, 3 \pmod{6}$. The matrix achieving equality would be $[K_m^0 K_m^1 K_m^2 T_{m,q-2}]$. Let B^c denote the $(0,1)$ -complement of a matrix B . Note that the $v \times \frac{a+b}{3} \binom{v}{2}$ simple matrix $[T_{v,a} T_{v,b}^c]$ has no submatrix $(a+b+1) \cdot (\mathbf{1}_2 \mathbf{0}_0)$.

Theorem 1.5 *Let $q > 2$ be given. There exists a constant $M = M(q)$ so that for $m > M$,*

$$\text{forb}(m, q \cdot (\mathbf{1}_2 \mathbf{0}_1)) = \left[\overbrace{\begin{matrix} 1 & 1 & \cdots & 1 \\ 1 & 1 & \cdots & 1 \\ 0 & 0 & \cdots & 0 \end{matrix}}^q \right] \leq m + 2 + \frac{q+1}{3} \binom{m}{2} \quad (6)$$

with equality for $m \equiv 1, 3 \pmod{6}$. If A is an $m \times \text{forb}(m, q \cdot (\mathbf{1}_2 \mathbf{0}_2))$ simple matrix with $m > M$ and $m \equiv 1, 3 \pmod{6}$, then A consists of all possible columns of sum 0, 1, 2, m and the columns of column sum 3 correspond to a simple triple system $T_{m,q-2}$ and A has no further columns. ■

Theorem 1.6 *Let $q > 2$ be given. There exists a constant $M = M(q)$ so that for $m > M$,*

$$\text{forb}(m, q \cdot (\mathbf{1}_2 \mathbf{0}_2)) = \left[\overbrace{\begin{matrix} 1 & 1 & \cdots & 1 \\ 1 & 1 & \cdots & 1 \\ 0 & 0 & \cdots & 0 \\ 0 & 0 & \cdots & 0 \end{matrix}}^q \right] \leq 2 + 2m + \frac{q+3}{3} \binom{m}{2}, \quad (7)$$

with equality for $m \equiv 1, 3 \pmod{6}$. If A is an $m \times \text{forb}(m, q \cdot (\mathbf{1}_2 \mathbf{0}_2))$ simple matrix with $m > M$ and $m \equiv 1, 3 \pmod{6}$, then there exist positive integers a, b with $a+b = q-3$ so that A consists of all possible columns of sum 0, 1, 2, $m-2$, $m-1$, m and the columns

of column sum 3 correspond to a simple triple system $T_{m,a}$ and the columns of column sum $m - 3$ correspond to the complement of a simple triple system $T_{m,b}$ and A has no further columns. ■

Thus the constructions for equality in Theorem 1.5 are $A = [K_m^0 K_m^1 K_m^2 T_{m,q-2} K_m^m]$ and the constructions for equality in Theorem 1.6 are found by selecting a, b positive integers with $a + b = q - 3$ and using $A = [K_m^0 K_m^1 K_m^2 T_{m,a} T_{m,b}^c K_m^{m-2} K_m^{m-1} K_m^m]$. For $m = q + 1$, the construction $A = [K_m^0 K_m^1 K_m^2 K_m^3 K_m^m]$ avoids $q \cdot (\mathbf{1}_2 \mathbf{0}_1)$ and exceeds the bound (8) and for $m = q + 1$, the construction $A = [K_m^0 K_m^1 K_m^2 K_m^3 K_m^{m-2} K_m^{m-1} K_m^m]$ and avoids $q \cdot (\mathbf{1}_2 \mathbf{0}_2)$ and exceeds the bound (9) so our theorems need some condition on m .

To prove Theorem 1.1 and Theorem 1.5, we prove the following:

Proposition 1.7 *Let $m, q > 2$ be given. Let A be an $m \times n$ $(0,1)$ -matrix so that no column of sum $0, 1, 2$, or m is repeated. Assume A has no configuration $q \cdot (\mathbf{1}_2 \mathbf{0}_1)$. Then there exists a constant M so that for $m > M$,*

$$n \leq m + 2 + \frac{q+1}{3} \binom{m}{2} \quad (8)$$

with equality for $m \equiv 1, 3 \pmod{6}$. If A is an $m \times \text{forb}(m, q \cdot (\mathbf{1}_2 \mathbf{0}_1))$ matrix with $m > M$ and $m \equiv 1, 3 \pmod{6}$, then A consists of all possible columns of sum $0, 1, 2, m$ once each and the columns of column sum 3 correspond to the triples of a 2-design $S_{q-2}(2, 3, m)$ and A has no further columns. ■

We see that Theorem 1.1 follows by taking a matrix A of column sums in $\{3, 4, \dots, m-1\}$ and with no configuration $(\lambda+1) \cdot (\mathbf{1}_2 \mathbf{0}_1)$ and adding the $\binom{m}{2} + m + 2$ columns of column sum $0, 1, 2$ and m to obtain a matrix A' . Now A' has no configuration $(\lambda+2) \cdot (\mathbf{1}_2 \mathbf{0}_1)$ and satisfies the hypotheses of Proposition 1.7 with $q = \lambda + 2$. Applying Proposition 1.7 yields Theorem 1.1. The bound of Theorem 1.5 follows directly from Proposition 1.7. To prove Theorem 1.2 and Theorem 1.6 we prove the following:

Proposition 1.8 *Let $m, q > 2$ be given. Let A be an $m \times n$ $(0,1)$ -matrix so that no column of sum $0, 1, 2, m - 2, m - 1$ or m is repeated. Assume A has no configuration $q \cdot (\mathbf{1}_2 \mathbf{0}_2)$. Then there exists a constant M so that for $m > M$,*

$$n \leq 2m + 2 + \frac{q+3}{3} \binom{m}{2} \quad (9)$$

with equality for $m \equiv 1, 3 \pmod{6}$. If A is an $m \times \text{forb}(m, q \cdot (\mathbf{1}_2 \mathbf{0}_2))$ matrix with $m > M$ and $m \equiv 1, 3 \pmod{6}$, then A consists of all possible columns of sum $0, 1, 2, m - 2, m - 1$ and m once each and there are two positive integers a, b satisfying $a + b = q - 3$ with the columns of column sum 3 correspond to the triples of a 2-design $S_a(2, 3, m)$ and the columns of column sum $m - 3$ correspond to the complements in $[m]$ of the blocks of a 2-design $S_b(2, 3, m)$ and A has no further columns. ■

We see that Theorem 1.2 follows by taking a matrix A of column sums in $\{3, 4, \dots, m-3\}$ and with no configuration $(\lambda+1) \cdot (\mathbf{1}_2 \mathbf{0}_2)$ and adding the $2\binom{m}{2} + 2m + 2$ columns of column sum $0, 1, 2, m-2, m-1$ and m to obtain a matrix A' . Now A' has no configuration $(\lambda+3) \cdot (\mathbf{1}_2 \mathbf{0}_2)$ and satisfies the hypotheses of Proposition 1.8 with $q = \lambda + 3$. Applying Proposition 1.8 yields Theorem 1.2. The bound of Theorem 1.6 follows directly from Proposition 1.8.

We could give a simpler direct proof of Theorem 1.1 by using the proof of Proposition 1.8 and ignoring certain column sums. We were originally motivated by the forbidden configuration bounds of Theorems 1.5 and Theorem 1.6.

The proofs of Proposition 1.7 and Proposition 1.8 use Turán's bound for the maximum number of edges in a graph with no complete graph of a certain size. We do not explicitly give values for M since the values as given by the proofs are unlikely to be of value but our proof shows we may take M to be $O(q^3)$. Proposition 1.7 for $q \cdot (\mathbf{1}_2 \mathbf{0}_1)$ is proven in Section 3 and Proposition 1.8 for $q \cdot (\mathbf{1}_2 \mathbf{0}_2)$ is proven in Section 4. The proofs are organized to highlight analogies with the proof of Theorem 1.3 but the details are different. We were surprised that exact bounds were obtained. We do not see how to extend our exact proofs to $F = t \cdot (\mathbf{1}_k \mathbf{0}_k)$ with $k \geq 3$ and moreover do not have the analogue of Dehon's lovely Theorem 1.4 to provide a construction of simple k -designs.

2 Exact Bound for $q \cdot (\mathbf{1}_1 \mathbf{0}_1)$

This section gives the proof of Theorem 1.3. We have broken it into lemmas. Assume A is a simple m -rowed matrix with no configuration $q \cdot (\mathbf{1}_1 \mathbf{0}_1)$. Let a_i denote the number of columns with either exactly i 1's or i 0's for $i = 0, 1, 2$ and let a_3 be the number of remaining columns. Without loss of generality, we may assume $a_0 = 2$ since the column of all 0's and the column of all 1's cannot contribute to $q \cdot (\mathbf{1}_1 \mathbf{0}_1)$. Thus $2 + a_1 + a_2 + a_3$ is the number of columns of A .

In [1], we establish that

$$2 + a_1 + a_2 + a_3 \leq \left\lfloor \frac{(q+1)m}{2} + \frac{(q-3)m}{2(m-2)} \right\rfloor + 2$$

and as noted in the Introduction, we can achieve equality for some small m . We wish to show that these small values of m are exceptional. We assume

$$a_1 + a_2 + a_3 > \left\lfloor \frac{(q+1)m}{2} \right\rfloor \tag{10}$$

and seek a contradiction.

Lemma 2.1 *Let A be an $m \times n$ simple matrix with no $q \cdot (\mathbf{1}_1 \mathbf{0}_1)$. Assume $m \geq 6$. Then*

$$(m-1)a_1 + 2(m-2)a_2 + 3(m-3)a_3 \leq (2q-2) \binom{m}{2} = (q-1)m(m-1). \tag{11}$$

Assume $n > \frac{q+1}{2}m + 2$. Then

$$2m - \frac{m(q-3)}{m-3} < a_1 \leq 2m, \quad (12)$$

$$a_3 < \frac{m(q-3)}{m-5}. \quad (13)$$

Proof: A column of k 1's contains $\binom{k}{1} \binom{m-k}{1}$ configurations $\mathbf{1}_1 \mathbf{0}_1$. Note that $\binom{k}{1} \binom{m-k}{1} \geq \binom{3}{1} \binom{m-3}{1}$ for $3 \leq k \leq m-3$. By the pigeonhole argument, there are at most $(2q-2) \binom{m}{2}$ configurations $\mathbf{1}_1 \mathbf{0}_1$ in A else there will be $2q-1$ in one of the $\binom{m}{2}$ pairs of rows and hence at least q with the 1 of the $\mathbf{1}_1 \mathbf{0}_1$ in the same row yielding the configuration $q \cdot (\mathbf{1}_1 \mathbf{0}_1)$. This yields (11). Given $m \geq 6$, we have $m-1 < 2(m-2) < 3(m-3)$. Substituting in (11),

$$a_1(m-1) + 2(m-2)(a_2 + a_3) \leq m(m-1)(q-1).$$

Using $a_2 + a_3 > \frac{q+1}{2}m - a_1$ from (10) we have

$$a_1(m-1) + 2(m-2)\left(\frac{q+1}{2}m - a_1\right) < m(m-1)(q-1)$$

and so

$$2m^2 - mq - 3m < (m-3)a_1$$

from which we deduce the lower bound of (12). The upper bound of (12) follows from counting all possible columns.

To show a_3 is small, use (11) to obtain

$$a_1(m-1) + 2(m-2)\left(\frac{q+1}{2}m - a_1 - a_3\right) + 3(m-3)a_3 < m(m-1)(q-1).$$

Rearranging yields

$$(m-5)a_3 < m(q-2m+3) + (m-3)a_1.$$

Substituting $a_1 \leq 2m$, we obtain (13). \blacksquare

Form two graphs G_0, G_1 from the columns of A where the vertex set for both graphs corresponds to the rows of A . We form a graph G_0 from the columns of A of column sum $m-2$ so that if there is a column of A with $m-2$ 1's and two 0's on rows i, j we add an edge (i, j) to G_0 . Similarly we form a graph G_1 from the columns of A of column sum 2 so that if there is a column of A with $m-2$ 0's and two 1's on rows i, j , then G_1 has the edge (i, j) . Define $d_0(i)$ and $d_1(i)$ to be the degrees of i in G_0 and G_1 respectively. Hence

$$a_2 = \frac{1}{2} \sum_{i=1}^m (d_0(i) + d_1(i)). \quad (14)$$

Using (10), we obtain

$$a_1 + \frac{1}{2} \sum_{i=1}^m (d_0(i) + d_1(i)) + a_3 > \frac{q+1}{2}m$$

Multiplying by 2 and substituting the upper bounds (12) for a_1 and (13) for a_3 , yields

$$\begin{aligned} \sum_{i=1}^m (d_0(i) + d_1(i)) &> (q+1)m - 4m - \frac{2m(q-3)}{m-5} \\ &= m(q-3) \left(1 - \frac{2}{m-5}\right). \end{aligned} \quad (15)$$

Thus the average value of $d_0(i) + d_1(i)$ is close to $q-3$.

The possible columns of column sum 1 or $m-1$ are as follows. Define e_i to be the m -rowed column with a 1 in row i and 0's elsewhere and let e_i^c be the (0,1)-complement of e_i . Define

$$E_1 = \{i : 1 \leq i \leq m \text{ and } e_i \text{ is not in } A\},$$

$$E_0 = \{i : 1 \leq i \leq m \text{ and } e_i^c \text{ is not in } A\}.$$

We have $a_1 = 2m - |E_0| - |E_1|$ and so $|E_1| + |E_0| < \frac{m(q-3)}{m-3}$ by (12). For convenience of counting define

$$\epsilon(i) = \begin{cases} 0 & \text{if } i \notin E_1 \cup E_0 \\ 1 & \text{if } i \in E_1 \setminus E_0 \text{ or } i \in E_0 \setminus E_1 \\ 2 & \text{if } i \in E_1 \cap E_0 \end{cases}. \quad (16)$$

Thus $\sum_{i=1}^m \epsilon(i) = |E_0| + |E_1|$.

Lemma 2.2 *Assume $m > 3q + 2$. Then for all $i = 1, 2, \dots, m$, we have $d_0(i) + d_1(i) \leq q - 3 + \epsilon(i)$.*

Proof: Assume the contrary that k is an index with $l = d_0(k) + d_1(k) \geq q - 2 + \epsilon(k)$. Let N_1 be the vertices/rows connected to k by no edges in either G_0 or G_1 . Let N_2 be the number of vertices connected to k by an edge in G_0 or an edge in G_1 but not both. Let N_3 be the number of vertices connected to k by edges in both G_0 and G_1 . We have

$$|N_1| + |N_2| + |N_3| = m - 1, \quad |N_2| + 2|N_3| = d_0(k) + d_1(k) = l. \quad (17)$$

Consider a row $i \neq k$. There are at most $2q - 2$ configurations $\mathbf{1}_1 \mathbf{0}_1$ contained in rows k, i of A and there are $4 - \epsilon(i) - \epsilon(k)$ configurations $\mathbf{1}_1 \mathbf{0}_1$ contained in rows k, i of A in the columns of column sum 1 or $m-1$ (corresponding to those columns e_k, e_k^c, e_i, e_i^c which are present in A). If $i \in N_1$ then each edge incident with either k or i in either G_0 or G_1 corresponds to a column of A of column sum 2 or $m-2$ that has the configuration

$\mathbf{1}_1\mathbf{0}_1$ in rows k, i and hence we have $d_1(k) + d_0(k) + d_1(i) + d_0(i)$ configurations $\mathbf{1}_1\mathbf{0}_1$ in these columns. Thus $d_1(k) + d_0(k) + d_1(i) + d_0(i) + (4 - \epsilon(i) - \epsilon(k)) \leq 2q - 2$ which yields

$$d_1(i) + d_0(i) \leq 2q - 6 - l + \epsilon(i) + \epsilon(k).$$

In the case $i \in N_2$ then we note that an edge in say G_0 joining k, i contributes 2 to $d_0(i) + d_0(k)$ but the corresponding column does not contain the configuration $\mathbf{1}_1\mathbf{0}_1$ in rows i, k . A similar argument holds for an edge (k, i) in G_1 . By the above analysis we obtain

$$d_1(i) + d_0(i) \leq 2q - 4 - l + \epsilon(i) + \epsilon(k).$$

In the case $i \in N_3$ then we note that the two edges in G_0 and G_1 joining k, i contributes 4 to $d_0(i) + d_1(i) + d_0(k) + d_1(k)$ but correspond to only two columns neither of which contain the configuration $\mathbf{1}_1\mathbf{0}_1$. By the above analysis we obtain

$$d_1(i) + d_0(i) \leq 2q - 2 - l + \epsilon(i) + \epsilon(k).$$

Summarizing, we have for $i \in N_j$ and $j = 1, 2, 3$ that

$$d_1(i) + d_0(i) \leq 2q - 6 + 2(j - 1) - l + \epsilon(i) + \epsilon(k). \quad (18)$$

Now we sum our upper bounds on $d_0(i) + d_1(i)$ over all rows $i \in [m] = \{k\} \cup N_1 \cup N_2 \cup N_3$ and use (15) to obtain

$$\begin{aligned} l + \sum_{j \in \{1, 2, 3\}} \sum_{i \in N_j} (2q - 6 + 2(j - 1) - l + \epsilon(i) + \epsilon(k)) \\ \geq \sum_{i=1}^m (d_0(i) + d_1(i)) > m(q - 3) \left(1 - \frac{2}{m - 5}\right) \end{aligned}$$

This simplifies to

$$\begin{aligned} l + (2q - 6)(|N_1| + |N_2| + |N_3|) + 2(|N_2| + 2|N_3|) - (m - 1)l + \\ + (|E_0| + |E_1| - \epsilon(k)) + (m - 1)\epsilon(k) > m(q - 3) \left(1 - \frac{2}{m - 5}\right) \end{aligned}$$

Using $|N_1| + |N_2| + |N_3| = m - 1$, $l = |N_2| + 2|N_3|$ from (17), and $|E_0| + |E_1| \leq \frac{m(q-3)}{m-3}$ and rearranging yields

$$(2q - 6)(m - 1) - (m - 4)l + (m - 4)\epsilon(k) + 2\epsilon(k) + \frac{m(q - 3)}{m - 3} > m(q - 3) - \frac{2m(q - 3)}{m - 5}$$

Using $-l + \epsilon(k) \leq -(q - 2)$ and $\epsilon(k) \leq 2$ and rearranging we get

$$\frac{m(q - 3)}{m - 3} + \frac{2m(q - 3)}{m - 5} > m - 2. \quad (19)$$

We can rewrite (19) as $0 > m^3 - (3q + 2)m^2 + (11q - 2)m - 30$ which is impossible for $m > 3q + 2$. This contradiction establishes the lemma. ■

Let

$$Y = \{i : d_0(i) + d_1(i) = q - 3 \text{ and } \epsilon(i) = 0\}$$

Lemma 2.3 *Assume $m > \max\{3q + 2, 8q - 19\}$. Then we may assume $|Y| \geq m/2$.*

Proof: We consider $[m]$ divided into Y , $E_0 \cup E_1$, and $[m] \setminus (Y \cup E_0 \cup E_1)$. We use Lemma 2.2. We have

$$\sum_{i \in E_0 \cup E_1} d_0(i) + d_1(i) \leq \sum_{i \in E_0 \cup E_1} (q - 3) + |E_0| + |E_1| = |E_0 \cup E_1|(q - 3) + |E_0| + |E_1|$$

using $\sum_{i=1}^m \epsilon(i) = |E_0| + |E_1|$. We readily compute $\sum_{i \in Y} d_0(i) + d_1(i) = |Y|(q - 3)$ and

$$\sum_{i \in [m] \setminus (Y \cup E_0 \cup E_1)} d_0(i) + d_1(i) \leq \sum_{i \in [m] \setminus (Y \cup E_0 \cup E_1)} (q - 4) \leq (m - |Y| - |E_0 \cup E_1|)(q - 4)$$

Summing we obtain

$$\sum_{i \in [m]} d_0(i) + d_1(i) \leq m(q - 3) + |E_0| + |E_1| - m + |Y| + |E_0 \cup E_1|$$

Now using (15), we deduce

$$|E_0| + |E_1| + |E_0 \cup E_1| + \frac{2m(q - 3)}{m - 5} > m - |Y|$$

We use $|E_0 \cup E_1| \leq |E_0| + |E_1| < \frac{2m(q-3)}{m-3}$ by (12) to obtain $\frac{2m(q-3)}{m-3} + \frac{2m(q-3)}{m-5} > m - |Y|$. Now for $m > 8q - 19$ (so that $m - 3 > m - 5 \geq 8(q - 3)$), we have $\frac{2m(q-3)}{m-3} + \frac{2m(q-3)}{m-5} \leq m/2$. Thus for $m > 8q - 19$, we may assume $|Y| \geq m/2$. ■

Let A_3 denote the submatrix of A formed by the columns of sum $3, 4, \dots, m - 3$. Then A_3 has a_3 columns. Let $A_3(Y)$ denote the submatrix of A_3 indexed by the rows of Y .

Lemma 2.4 *Assume $m > \max\{3q + 2, 8q - 19\}$. Then $A_3(Y)$ has no configuration $\mathbf{1}_1 \mathbf{0}_1$.*

Proof: Assume there is a column α in A_3 which has both 0's and 1's in the rows indexed by Y . By taking the (0,1)-complement of A if necessary, we may assume the number of 1's in those rows is at least $|Y|/2 \geq m/4$. Consider a row $i \in Y$ where α has a 0. Then there exists a row $j \in Y$ where α has a 1 such that rows i, j are not connected in G_0 or G_1 , since i is connected to at most $q - 3$ rows and $|Y|/2 \geq m/4 > q - 3$. Given $i, j \in Y$, we have $d_0(i) + d_1(i) = d_0(j) + d_1(j) = q - 3$ and $\epsilon(i) = \epsilon(j) = 0$. Given that i, j are not

connected in G_0 or G_1 , we have $2(q-3)$ copies of the configuration $\mathbf{1}_1\mathbf{0}_1$ on rows i, j in the columns (of A) of exactly two 1's or exactly two 0's. Given $Y \cap (E_1 \cup E_2) = \emptyset$, we have 4 copies of the configuration $\mathbf{1}_1\mathbf{0}_1$ in the columns (of A) of one 1 or one 0 and in rows i, j . But α has $\mathbf{1}_1\mathbf{0}_1$ in rows i, j and so we find $q \cdot (\mathbf{1}_1\mathbf{0}_1)$ in A , a contradiction. This establishes the lemma. \blacksquare

Proof of Theorem 1.3: We obtain a contradiction from assuming (10) and $m > \max\{3q+2, 8q-19\}$ and thus establish (5). By Lemma 2.4, each column of A_3 has either all 1's or all 0's on the rows of Y . Considering column sums, every column in A_3 which has all 1's on rows Y , has at least three 0's and every column in A_3 which has all 0's on rows Y , has at least three 1's. For $i \in [m] \setminus Y$, let $t_0(i)$ denote the number of 0's in columns of column sum in $\{3, 4, \dots, m-3\}$ which are all 1's on Y and let $t_1(i)$ denote the number of 1's in columns of column sum in $\{3, 4, \dots, m-3\}$ which are all 0's on Y . Counting yields

$$\sum_{i \in [m] \setminus Y} (t_0(i) + t_1(i)) \geq 3a_3. \quad (20)$$

Let $i \in [m] \setminus Y$ be given. We wish to establish

$$d_0(i) + d_1(i) \leq q - 3 + \epsilon(i) - t_0(i) - t_1(i) \quad (21)$$

We use a similar argument as Lemma 2.2. Consider a column of A_3 which is all 0's on rows of Y . Then the column has a configuration $\mathbf{1}_1\mathbf{0}_1$ in rows i, k for any choice of $k \in Y$. A similar remarks holds for columns of A_3 which are all 1's on rows of Y . Let X denote all the neighbours of i in G_0 and in G_1 . We have $|X| \leq d_0(i) + d_1(i) \leq q-1$ using Lemma 2.2. Given $|Y| > m/2 > q-1$, we can select a $k \in Y$ with $k \notin X$. Now the columns of sum 1 or $m-1$ in A have $4 - \epsilon(i) - \epsilon(k) = 4 - \epsilon(i)$ configurations $\mathbf{1}_1\mathbf{0}_1$ in rows i, k (since $\epsilon(k) = 0$). Given that $k \notin X$, the columns of column sum 2 or $m-2$ have $d_0(i) + d_1(i) + d_0(k) + d_1(k) = d_0(i) + d_1(i) + q-3$ configurations $\mathbf{1}_1\mathbf{0}_1$ in rows i, k . The columns of sum at least 3 and at most $m-3$ have at least $t_0(i) + t_1(i)$ configurations $\mathbf{1}_1\mathbf{0}_1$ in rows i, k for that choice of k . Rows i, k of A have at most $2(q-1)$ such configurations and so we obtain (21).

Combining twice (10) and (14) we have

$$2a_1 + \sum_{i=1}^m (d_0(i) + d_1(i)) + 2a_3 > m(q+1).$$

Using $a_1 = 2m - |E_0| + |E_1|$, substituting $d_0(i) + d_1(i) = q-3$ for $i \in Y$ and using (21),

$$\sum_{i \in Y} (q-3) + \sum_{i \in [m] \setminus Y} (q-3 + \epsilon(i) - t_0(i) - t_1(i)) + 2a_3 > m(q+1) - 2(2m - |E_0| + |E_1|).$$

Now using (20) and $\sum_{i \in [m] \setminus Y} \epsilon(i) \leq |E_0| + |E_1|$,

$$|Y|(q-3) + (m-|Y|)(q-3) + |E_0| + |E_1| - a_3 > m(q-3) + 2(|E_0| + |E_1|)$$

which yields the contradiction (even for $a_3 = 0$ and $|E_0| + |E_1| = 0$)

$$-a_3 > |E_0| + |E_1|.$$

This final contradiction establishes (5). \blacksquare

One could note that for a matrix A to achieve equality, we would have $a_3 = 0$ and $|E_0| + |E_1| = 0$ and so $a_1 = 2m$. This suggests that A would have to correspond to the construction given in the Introduction or its (0,1)-complement.

3 Exact Bound for $q \cdot (\mathbf{1}_2 \mathbf{0}_1)$

We are able to generalize the argument for Theorem 1.3 following a similar series of Lemmas to obtain Proposition 1.7. We do not explicitly calculate the smallest possible constant M for our proof (following the argument yields that M is $O(q^3)$), believing that our argument does not give a realistic values for M . Let A be an $m \times n$ (0,1)-matrix with no configuration $q \cdot (\mathbf{1}_2 \mathbf{0}_1)$ so that there are no repeated columns of sum 0, 1, 2, m . We wish to ignore the $m+2$ possible columns of sum 0, 1, m since they cannot contribute to a configuration $q \cdot (\mathbf{1}_2 \mathbf{0}_1)$. So assume A has column sums between 2 and $m-1$, inclusive. Assume $n > \frac{q+1}{3} \binom{m}{2}$. We wish to arrive at a contradiction to prove (8).

For $i = 2, 3$, let a_i denote the number of columns of column sum i in A and let a_4 denote the number of columns of column sum at least 4 in A . Note that the definition of a_1, a_2, \dots is different in this section from Sections 2 and 4. Note that we do not allow repeated columns of sum 2. We have by assumption that

$$a_2 + a_3 + a_4 > \frac{q+1}{3} \binom{m}{2}. \quad (22)$$

Lemma 3.1 *Let m, q be given. Let A be an $m \times n$ simple matrix with no $q \cdot (\mathbf{1}_2 \mathbf{0}_1)$. Assume $m \geq 6$ and (22). Then*

$$\binom{2}{2} \binom{m-2}{1} a_2 + \binom{3}{2} \binom{m-3}{1} a_3 + \binom{4}{2} \binom{m-4}{1} a_4 \leq \binom{m}{3} 3(q-1). \quad (23)$$

There exists positive constants c_1, c_2 so that

$$\binom{m}{2} - c_1 m \leq a_2 \leq \binom{m}{2}, \quad (24)$$

$$a_4 \leq c_2 m. \quad (25)$$

Proof: We note that a column of column sum k has $\binom{k}{2} \binom{m-k}{1}$ configurations $\mathbf{1}_2 \mathbf{0}_1$ and note that $\binom{k}{2} \binom{m-k}{1} \geq \binom{4}{2} \binom{m-4}{1}$ for $4 \leq k \leq m-1$. Counting the configurations $\mathbf{1}_2 \mathbf{0}_1$ and using the pigeonhole argument yields (23)

For $m \geq 6$ we have $\binom{3}{2} \binom{m-3}{1} < \binom{4}{2} \binom{m-4}{1}$. Hence

$$(m-2)a_2 + 3(m-3)(a_3 + a_4) \leq \binom{m}{3} 3(q-1)$$

From (22), we have $a_3 + a_4 \geq \frac{q+1}{3} \binom{m}{2} - a_2$. We substitute and obtain

$$(m-3) \binom{m}{2} (q+1) - \binom{m}{3} 3(q-1) \leq \left(3(m-3) - (m-2) \right) a_2$$

which simplifies as

$$\binom{m}{2} (2m - q - 5) \leq (2m - 7) a_2$$

from which we deduce that there is a constant c_1 (will depend on q) so that first half of (24) holds. The second half of (24) follows from the fact that no column of sum 2 is repeated.

In a similar way we have

$$(m-2)a_2 + 3(m-3) \left(\frac{q+1}{3} \binom{m}{2} - a_2 - a_4 \right) + 6(m-4)a_4 \leq \binom{m}{3} 3(q-1)$$

and when we substitute the upper bound of (24), we deduce that there is a constant c_2 (will depend on q) so that (25) holds. \blacksquare

Partition A into three parts: A_2 consists of the columns of column sum 2, A_3 is the columns of column sum 3 and A_4 is the columns of column sum greater or equal than 4. We will refer to $\mathcal{A}_2, \mathcal{A}_3$ using the notations of (3) and (4). Note that \mathcal{A}_3 is a multiset and \mathcal{A}_2 is a set given that there are no repeated columns of sum 2. Considering the columns of column sum 2, we adapt $\epsilon(i)$ of (16). Note that for convenience we represent every pair $\{i, j\}$ by ij and so $ij \equiv ji$. We are not interested in ordered pairs in this context. Define

$$\epsilon(ij) = \begin{cases} 1 & \text{if } \{i, j\} \notin \mathcal{A}_2 \\ 0 & \text{if } \{i, j\} \in \mathcal{A}_2 \end{cases}, \quad E = \{ij : \epsilon(ij) = 1\} \quad .$$

Thus

$$a_2 = \binom{m}{2} - \sum_{ij} \epsilon(ij) = \binom{m}{2} - |E|. \quad (26)$$

We deduce from (24) that $|E| \leq c_1 m$.

We adapt the definitions of the degrees d_0, d_1 of Section 2 by using a hypergraph degree definitions applied to the multiset $\mathcal{A}_3 = \{B_1, B_2, \dots\}$. Define

$$d(ij) = |\{s : B_s \in \mathcal{A}_3 \text{ and } i, j \in B_s\}|$$

Then

$$3a_3 = \sum_{\{i,j\} \subseteq [m]} d(ij). \quad (27)$$

Let

$$\mathcal{U}(pt) = \{r : \{p, t, r\} \in \mathcal{A}_3\}.$$

Since $m > q + 2$ and we are avoiding $q \cdot (\mathbf{1}_2 \mathbf{0}_1)$ in A_3 then $|\mathcal{U}(pt)| < q$. Also let

$$\mathcal{T}(r) = \{pt : \{p, t, r\} \in \mathcal{A}_3\}.$$

Since for every $x \in [m]$ with $x \neq r$, $|\mathcal{U}(rx)| < q$ we have $|\mathcal{T}(r)| < \frac{(m-1)q}{2}$. Note that $\mathcal{U}(pt)$ and $\mathcal{T}(r)$ are the generalizations of X (found after (21)) given in the proof of Theorem 1.3.

Lemma 3.2 *We have*

$$d(ij) \leq (q - 2) + \epsilon(ij). \quad (28)$$

Proof: Since $m > q + 2 \geq |\mathcal{U}(ij)| + 2$, for every pair ij , we can find row $k \neq i, j$ so that $k \notin \mathcal{U}(ij)$. Now the number of submatrices

$$\begin{matrix} i \\ j \\ k \end{matrix} \begin{bmatrix} 1 \\ 1 \\ 0 \end{bmatrix} \quad (29)$$

in A_3 is $d(ij)$ (since $d(ij)$ is the number of triples i, j, l corresponding to columns in A_3 and each such column yields the submatrix since $k \notin \mathcal{U}(ij)$) and the number of submatrices (29) in A_2 is $1 - \epsilon(ij)$. Thus

$$d(ij) + 1 - \epsilon(ij) \leq q - 1$$

and hence (28) holds. \blacksquare

Let

$$Y = \{ij : d(ij) = q - 2 \text{ and } \epsilon(ij) = 0\}$$

Lemma 3.3 *There exists a constant c_3 so that*

$$|Y| \geq \binom{m}{2} - c_3 m \quad (30)$$

Proof: We partition the $\binom{m}{2}$ pairs ij into 3 parts: Y , E and the rest. By Lemma 3.2, for each $ij \in E$ we have $d(ij) \leq (q - 2) + 1$. Note that for $ij \notin Y \cup E$, we have $\epsilon(ij) = 0$ and so $d(ij) \leq (q - 2) - 1$ else $ij \in Y$. Thus from (27)

$$3a_3 = \sum_{ij} d(ij) \leq \left((q - 2)|Y| + ((q - 2) + 1)|E| + ((q - 2) - 1) \left(\binom{m}{2} - |Y| - |E| \right) \right)$$

Hence

$$a_3 \leq \frac{1}{3} \left((q - 2) \binom{m}{2} + |E| - \binom{m}{2} + |Y| + |E| \right) \quad (31)$$

Substituting estimates of a_2, a_3, a_4 from (26), (31), (25) into (22), we have

$$\binom{m}{2} - |E| + \frac{1}{3} \left((q-2) \binom{m}{2} + 2|E| - \binom{m}{2} + |Y| \right) + c_2 m > \frac{q+1}{3} \binom{m}{2}$$

We deduce $-\frac{1}{3}|E| + \frac{1}{3}|Y| + c_2 m > \frac{1}{3} \binom{m}{2}$ and so there exists a constant $c_3 = 3c_2$ so that (30) holds. ■

Form a graph G of m vertices corresponding to the rows of A and with edges (i, j) if and only if $ij \in Y$. Thus by Lemma 3.3, the number of edges of G is at least $\binom{m}{2} - c_3 m$. By Turán's Theorem [7], a graph with more than $\frac{m^2}{2} - \frac{m^2}{2(k-1)}$ edges has a clique of k vertices. Thus G has large cliques. Let c_4 be a constant chosen so that for any choices of i, j, k the following three inequalities hold.

$$\begin{aligned} \binom{c_4 \sqrt{m}/2}{2} &> \frac{m-1}{2} q (> |\mathcal{T}(k)|), & \frac{c_4 \sqrt{m}}{2} &> q (> |\mathcal{U}(ij)|), \\ \binom{c_4 \sqrt{m}}{2} &> \frac{m-1}{2} q + 3m (\geq |\mathcal{T}(k)| + |\mathcal{U}(ij)|) \end{aligned} \quad (32)$$

By Turán's argument, there exists an M so that for $m \geq M$, we can find a clique of $c_4 \sqrt{m}$ vertices in G . Let the vertices in this clique be denoted B . Thus for $i, j \in B$ we have $d(ij) = q-2$ and $\epsilon(ij) = 0$. Let $A_4(B)$ be the submatrix of A_4 of the rows indexed by B .

Lemma 3.4 *Assume $m > M$. Then $A_4(B)$ has no configuration $\mathbf{1}_2 \mathbf{0}_1$.*

Proof: Consider a column α of A_4 . We consider two cases based on whether there are more 1's or more 0's in the rows B . Assume α has at least $\frac{c_4 \sqrt{m}}{2}$ 1's in rows of B . Assume α has a 0 in row $k \in B$. Then by the first inequality (32), there is a pair $ij \notin \mathcal{T}(k)$ with $i, j \in B$. Thus there are $q-2$ columns of column sum 3 with the submatrix (29) using $d(ij) = q-2$ and 1 column of column sum 2 with the submatrix (29) using $\epsilon(ij) = 0$ and column α has 1 further submatrix (29) which creates the configuration $q \cdot (\mathbf{1}_2 \mathbf{0}_1)$, a contradiction. So α has no configuration $\mathbf{1}_2 \mathbf{0}_1$.

Assume α of A_4 that has at least $\frac{c_4 \sqrt{m}}{2}$ 0's in the rows of B . Assume α has 1's in rows $i, j \in B$. Then there is a row $k \in B$ where α has a 0 in row k and $k \notin \mathcal{U}(ij)$ by the second inequality of (32). For that choice of k and using $d(ij) = q-2$, there are $q-2$ columns of column sum 3 with the submatrix (29). There is one column of column sum 2 with the submatrix (29) using $\epsilon(ij) = 0$ and the α has one further submatrix (29) which creates the configuration $q \cdot (\mathbf{1}_2 \mathbf{0}_1)$, a contradiction. Thus α has no configuration $\mathbf{1}_2 \mathbf{0}_1$. ■

Lemma 3.5 *Assume $m > M$. Then the inequality (8) holds.*

Proof: We obtain a contradiction from assuming $m > M$ and (22) and thus establish (8). Our proof considers the a_4 columns of A_4 (which are the columns of column sum at least 4 and at most $m - 1$).

From Lemma 3.4, each column in A_4 either has at most one 1 or has no 0's in the rows of B . Let A_4^0 be those columns of A_4 with at most one 1 in the rows of B and hence at least three 1's in the rows $[m] \setminus B$. Let a_4^0 be the number of columns in A_4^0 . Let A_4^1 be those columns of A_4 with no 0's in the rows of B and hence at least one 0 in the rows $[m] \setminus B$. Let a_4^1 be the number of columns of A_4^1 . We have $a_4^0 + a_4^1 = a_4$.

For a pair ij with $i, j \in [m] \setminus B$, let $t(ij)$ count the number of columns of A_4^0 with 1's in both rows i and j . Each column with at most one 1 in B has at least three 1's in $[m] \setminus B$ and hence 1's in at least $\binom{3}{2} = 3$ pairs ij with $i, j \in [m] \setminus B$. We have verified that

$$\sum_{ij: i, j \in [m] \setminus B} t(ij) \geq 3a_4^0. \quad (33)$$

We must work harder to get an analog of (33) for A_4^1 . Assume $a_4^1 > 0$. For a pair ij with $i, j \in B$ and $k \in [m] \setminus B$ with $ij \notin \mathcal{T}(k)$, let $t(ij, k)$ denote the number of submatrices (29) in A_4^1 . When $ij \in \mathcal{T}(k)$, set $t(ij, k) = 0$. For a pair ij with $i, j \in B$, let

$$t(ij) = \max_{k \in [m] \setminus B} t(ij, k) \quad (34)$$

Each column α in A_4^1 has at least one row, say $l \in [m] \setminus B$ with a 0. For column α , we know $|\mathcal{T}(l)| < \frac{m-1}{2}q$ and at the same time there are $\binom{c_4\sqrt{m}}{2}$ pairs ij with $i, j \in B$ and so there are at least $\binom{c_4\sqrt{m}}{2} - \frac{m-1}{2}q$ pairs ij with $i, j \in B$ with $ij \notin \mathcal{T}(l)$. Thus by the third inequality of (32), column α contributes at least $3m$ to the sum $\sum_{ij \notin \mathcal{T}(l)} t(ij, l)$ and so

$$\sum_{l \in [m] \setminus B} \sum_{ij: i, j \in B} t(ij, l) > 3ma_4^1$$

Thus by (34),

$$(m - |B|) \cdot \sum_{ij: i, j \in B} t(ij) > \sum_{l \in [m] \setminus B} \sum_{ij: i, j \in B} t(ij, l)$$

and so we deduce that

$$\sum_{ij: i, j \in B} t(ij) > \frac{3ma_4^1}{m - |B|} > 3a_4^1. \quad (35)$$

For a pair ij with $i \in B$ and $j \in [m] \setminus B$ or vice versa, let $t(ij) = 0$. We add (33) and (35) together to get

$$\sum_{ij} t(ij) \geq 3a_4, \quad (36)$$

with strict inequality if $a_4^1 > 0$.

We are able to extend Lemma 3.2 and establish

$$d(ij) \leq q - 2 + \epsilon(ij) - t(ij) \quad (37)$$

By Lemma 3.2, we need only consider ij with $t(ij) > 0$. Given the definition of $t(ij)$, we need only consider the two cases: $i, j \in [m] \setminus B$ or $i, j \in B$.

In the former case we note that each of the $t(ij)$ columns of A_4^0 with 1's in both rows i and j have at most one 1 in rows of B . With $|B| > 2q$, (by the second inequality of (32)) we deduce that $t(ij) < q$ else we will find the configuration $q \cdot (\mathbf{1}_2 \mathbf{0}_1)$ in A_4^0 in the rows i, j and a row of B . Now in these $t(ij)$ columns of A_4^0 , at least $|B| - q + 1$ rows of B are all 0's. Again using the second inequality of (32) that $|B| > 2q$, we can find some $k \in B$ with $k \notin \mathcal{U}(ij)$ and all the $t(ij)$ columns have 0's in row k . Now there are $d(ij)$ submatrices (29) in columns of sum 3, $(1 - \epsilon(ij))$ submatrices (29) in columns of sum 2, and $t(ij)$ submatrices (29) in columns of sum 4 or more. The total is at most $q - 1$ since otherwise we would have the configuration $q \cdot (\mathbf{1}_2 \mathbf{0}_1)$ and this yields $d(ij) + (1 - \epsilon(ij)) + t(ij) \leq q - 1$. This is (37).

In the latter case with $i, j \in B$, we select k so that $t(ij, k) = t(ij)$. Thus $k \notin \mathcal{U}(ij)$ and also there are at least $t(ij)$ submatrices (29) in columns of A_4^1 . Thus we can now follow the same argument as in the former case to establish (37).

Now using (37) and (36),

$$3a_3 = \sum_{ij} d(ij) \leq \sum_{ij} \left(q - 2 + \epsilon(ij) - t(ij) \right) = (q - 2) \binom{m}{2} + |E| - 3a_4. \quad (38)$$

Substituting (26), (38), and (25) in (22) we obtain

$$\binom{m}{2} - |E| + \frac{1}{3} \left((q - 2) \binom{m}{2} + |E| - 3a_4 \right) + a_4 > \frac{q + 1}{3} \binom{m}{2}.$$

Simplifying and rearranging,

$$-\frac{2}{3}|E| > 0$$

which is a contradiction (even for $|E| = 0$) and this establishes (8). \blacksquare

Proof of Proposition 1.7: Lemma 3.5 establishes most of Proposition 1.7 but we are also interested in cases when the bound is achieved. Assume $m > M$ and $m \equiv 1, 3 \pmod{6}$. We now consider an m -rowed simple matrix A which has no configuration $q \cdot (\mathbf{1}_2 \mathbf{0}_1)$ and with $\binom{m}{0} + \binom{m}{1} + \frac{q+1}{3} \binom{m}{2} + \binom{m}{m}$ columns. One repeats the previous lemmas and arguments replacing the inequality (22) with the equation

$$a_2 + a_3 + a_4 = \frac{q + 1}{3} \binom{m}{2}. \quad (39)$$

We wish to show $a_2 = \binom{m}{2}$, $a_4 = 0$, $a_3 = \frac{q-2}{3} \binom{m}{2}$. Now Lemma 3.1 holds with (22) as an equality. We deduce the same bounds for $\mathcal{U}(rx)$ and $\mathcal{T}(r)$. Lemma 3.2 still holds

since the final contradiction does not require the strict inequality of (22) merely the equality of (39). Lemma 3.3 holds and we can choose B as large as possible but at least satisfying the three inequalities (32). Lemma 3.4 continues to hold.

We use (39) and following the argument of Lemma 3.5, we deduce that $E = \emptyset$ and so $a_2 = \binom{m}{2}$. Also we deduce that

$$\sum_{ij} t(ij) = 3a_4$$

and as a result of the strict inequality in (35), we can deduce that $a_4^1 = 0$.

Assume $a_4 = a_4^0 > 0$ and consider α in A_4 with column sum 4 and with 1's in rows i, j, k, l where $i \in B$ and $j, k, l \in \{1, 2, \dots, m\} \setminus B$. Choose $r \in B \setminus i$ then α has 1's in rows i, j and 0's in row r . Using $E = \emptyset$, we deduce that for this particular i, j we have $d(ij) \leq (q-2) - 1$. This yields a slight variant of (38):

$$3a_3 = \sum_{ij} d(ij) \leq \sum_{ij} ((q-2) + \epsilon(ij) - t(ij)) - 1.$$

The extra '-1' is sufficient to obtain a contradiction when we substitute for a_2, a_3, a_4 in (39). We then deduce $a_4 = 0$.

With $a_4 = 0$ and $a_2 = 2\binom{m}{2}$, we deduce $a_3 = \frac{q-3}{3}\binom{m}{2}$ using (39). Given that $\epsilon(ij) = 0$ for all ij and using Lemma 3.2, we deduce $d(ij) = q-2$ for all pairs ij and so $B = \{1, 2, \dots, m\}$. From this we can readily conclude that the columns of column sum 3 correspond to a 2-design $S_{q-2}(2, 3, m)$ and A has no further columns. ■

4 Exact Bound for $q \cdot (\mathbf{1}_2 \mathbf{0}_2)$

We generalize our proof of Proposition 1.7 given in Section 3 to prove Proposition 1.8. Again we do not explicitly calculate the smallest possible constant M but we note that we can take M to be $O(q^3)$.

Let A be a $m \times n$ matrix with no $q \cdot (\mathbf{1}_2 \mathbf{0}_2)$. Assume that there are no repeated columns of sums $0, 1, 2, m-2, m-1, m$. We will assume $n > 2 + 2m + \binom{m}{2} \frac{q+3}{3}$. Let a_i denote the number of columns with either exactly i 1's or i 0's for $i = 0, 1, 2, 3$ and let a_4 be the number of remaining columns. We may assume $a_0 = 2$ and $a_1 = 2m$ since all columns of column sum $0, 1, m-1$ or m do not contain the configuration $\mathbf{1}_2 \mathbf{0}_2$. Thus

$$a_2 + a_3 + a_4 > \binom{m}{2} \frac{q+3}{3}. \quad (40)$$

Lemma 4.1 *Assume A is an $m \times n$ simple matrix with no configuration $q \cdot (\mathbf{1}_2 \mathbf{0}_2)$ and (40) holds. Then there exists an m_0 so that for $m > m_0$,*

$$\binom{2}{2} \binom{m-2}{2} a_2 + \binom{3}{2} \binom{m-3}{2} a_3 + \binom{4}{2} \binom{m-4}{2} a_4 \leq 6 \binom{m}{4} (q-1). \quad (41)$$

Also there exist constants c_1, c_2 so that

$$2\binom{m}{2} - c_1 m \leq a_2 \leq 2\binom{m}{2} \quad (42)$$

$$a_4 \leq c_2 m \quad (43)$$

Proof: A column in A of column sum k has $\binom{k}{2}\binom{m-k}{2}$ configurations $\mathbf{1}_2\mathbf{0}_2$. Note that $\binom{k}{2}\binom{m-k}{2} \geq \binom{4}{2}\binom{m-4}{2}$ for $4 \leq k \leq m-4$. By the pigeonhole principle, there are at most $6(q-1)\binom{m}{2}$ configurations $\mathbf{1}_2\mathbf{0}_2$ in A . We obtain (41). There exist an m_0 , such that for $m > m_0$, $\binom{3}{2}\binom{m-3}{2} < \binom{4}{2}\binom{m-4}{2}$. Substituting in (41),

$$\binom{m-2}{2}a_2 + 3\binom{m-3}{2}(a_3 + a_4) \leq 6(q-1)\binom{m}{4}$$

which yields using $a_3 + a_4 > \binom{m}{2}\frac{q+3}{3} - a_2$ and rearranging

$$\binom{m-2}{2}a_2 + 3\binom{m-3}{2}\left[\binom{m}{2}\frac{q+3}{3} - a_2\right] < 6(q-1)\binom{m}{4}.$$

Therefore,

$$\binom{m-3}{2}\binom{m}{2}(q+3) - 6(q-1)\binom{m}{4} < \left(3\binom{m-3}{2} - \binom{m-2}{2}\right)a_2. \quad (44)$$

The leading term on the righthand side is exactly m^4 while the leading coefficient of a_2 on the lefthand side is exactly m^2 . Thus (44) implies that there exists some constant c_1 so that the lower bound of (42) holds. The upper bound of (42) follows from the fact no column of sum 2 or $m-2$ is repeated.

We can also bound a_4 . From (40), we have $a_3 > \frac{q+3}{3}\binom{m}{2} - a_2 - a_4$. Using (41) we have

$$\binom{m-2}{2}a_2 + 3\binom{m-3}{2}\left[\frac{q+3}{3}\binom{m}{2} - a_2 - a_4\right] + 6\binom{m-4}{2}a_4 \leq 6(q-1)\binom{m}{4}.$$

Then

$$\begin{aligned} & \left[6\binom{m-4}{2} - 3\binom{m-3}{2}\right]a_4 \\ & \leq 6(q-1)\binom{m}{4} - (q+3)\binom{m-3}{2}\binom{m}{2} + a_2\left[3\binom{m-3}{2} - \binom{m-2}{2}\right]. \end{aligned}$$

Substituting $a_2 \leq 2\binom{m}{2}$ and rearranging we have

$$\left[6\binom{m-4}{2} - 3\binom{m-3}{2}\right]a_4 \leq \frac{m(m-1)(m-3)}{4}(2q-6) \quad (45)$$

Then (45) implies that there exist some constant c_2 so that (43) holds. \blacksquare

We could have produced the bound $a_4 \leq (2q-6)\frac{m}{6} + c'_2$ for some constant c'_2 , but this is of little help. Now we form analogs of the degrees d_0, d_1 of Section 2 by defining A_3 as the submatrix of A of the columns of column sum 3 and defining A_{m-3} as the submatrix of A of the columns of column sum $m-3$. We refer to the mutisets $\mathcal{A}_3 = \{B_1, B_2, \dots\}$, $\mathcal{A}_{m-3} = \{C_1, C_2, \dots\}$ using the notations of (3) and (4). Define

$$d_1(ij) = |\{s : B_s \in \mathcal{A}_3 \text{ and } i, j \in B_s\}|, \quad d_0(ij) = |\{s : C_s \in \mathcal{A}_{m-3} \text{ and } i, j \notin C_s\}|$$

Recalling $a_3 = |\mathcal{A}_3| + |\mathcal{A}_{m-3}|$, we note

$$3a_3 = \sum_{\{i,j\} \subset [m]} (d_0(ij) + d_1(ij)) \quad (46)$$

Define e_{ij} to be the m -rowed column with 1 in rows i and j and 0's elsewhere, and let e_{ij}^c be the $(0,1)$ -complement of e_{ij} . These are the possible columns of column sum 2 or $m-2$. Define

$$E_1 = \{ij : \{i, j\} \subset [m] \text{ and } e_{ij} \text{ is not in } A\}$$

$$E_0 = \{ij : \{i, j\} \subset [m] \text{ and } e_{ij}^c \text{ is not in } A\}$$

For convenience of counting define

$$\epsilon(ij) = \begin{cases} 0 & \text{if } ij \notin E_1 \cup E_0 \\ 1 & \text{if } ij \in E_1 \setminus E_0 \text{ or } ij \in E_0 \setminus E_1 \\ 2 & \text{if } ij \in E_1 \cap E_0 \end{cases} \quad (47)$$

Thus

$$a_2 = 2 \binom{m}{2} - \sum_{i,j \subset [m]} \epsilon(ij) = 2 \binom{m}{2} - (|E_1| + |E_0|), \quad (48)$$

and given (42) we have $|E_1| + |E_0| \leq c_1 m$

We note for a quadruple of rows p, t, r, s that are at most $2q-2$

$$\text{submatrices } \begin{matrix} p \\ t \\ r \\ s \end{matrix} \begin{bmatrix} 1 \\ 1 \\ 0 \\ 0 \end{bmatrix} \text{ or submatrices } \begin{matrix} p \\ t \\ r \\ s \end{matrix} \begin{bmatrix} 0 \\ 0 \\ 1 \\ 1 \end{bmatrix} \quad (49)$$

else A has the configuration $q \cdot (\mathbf{1}_2 \mathbf{0}_2)$. For disjoint pairs pt and rs (i.e. $\{p, t\} \cap \{r, s\} = \emptyset$) we say *pair pt has triple overlapping rs* if and only if at least one of submatrices

$$\begin{matrix} p \\ t \\ r \\ s \end{matrix} \begin{bmatrix} 1 \\ 1 \\ 1 \\ 0 \end{bmatrix} \text{ or } \begin{matrix} p \\ t \\ r \\ s \end{matrix} \begin{bmatrix} 1 \\ 1 \\ 0 \\ 1 \end{bmatrix}$$

appears in columns of column sum 3 or at least one of submatrices

$$\begin{array}{c} p \\ t \\ r \\ s \end{array} \begin{bmatrix} 0 \\ 0 \\ 0 \\ 1 \end{bmatrix} \quad \text{or} \quad \begin{array}{c} p \\ t \\ r \\ s \end{array} \begin{bmatrix} 0 \\ 0 \\ 1 \\ 0 \end{bmatrix}$$

appears in columns of column sum $m - 3$. This definition is not symmetric in the pair pt, rs . Note that columns of three 1's that have 1's on rows p, t yet no 1's on rows r, s or vice versa have 1's on rows r, s yet no 1's on rows p, t contribute to (49). Similarly for columns with three 0's. Let

$$\mathcal{U}(pt) = \{ij : \{i, j\} \subset [m], \text{ pair } pt \text{ has triple overlapping } ij\},$$

$$\mathcal{T}(pt) = \{ij : \{i, j\} \subset [m], \text{ pair } ij \text{ has triple overlapping } pt\}.$$

Given $m > q + 2$, we cannot have the submatrix $q \cdot (\mathbf{1}_2 \mathbf{0}_0)$ in rows p, t in columns of column sum 3 else we would have the configuration $q \cdot (\mathbf{1}_2 \mathbf{0}_2)$ (and so there are at most $q - 1$ columns of column sum 3 with 1's in rows p, t). Similarly, we cannot have the submatrix $q \cdot (\mathbf{1}_0 \mathbf{0}_2)$ in rows p, t in columns of column sum $m - 3$. To bound $\mathcal{U}(pt)$, we note that $\binom{m-2-(q-1)}{2}$ counts the number of pairs ij disjoint from pt that avoids $q - 1$ further rows. Thus the number of pairs ij where pt overlaps ij using a column of column sum 3 is at most $\binom{m-2}{2} - \binom{m-2-(q-1)}{2}$. Similarly, the number of pairs ij where pt overlaps ij using a column of column sum $m - 3$ is at most $\binom{m-2}{2} - \binom{m-2-(q-1)}{2}$. Thus there exists a constant c_3 depending only on q so that

$$|\mathcal{U}(pt)| \leq 2 \left(\binom{m-2}{2} - \binom{m-2-(q-1)}{2} \right) \leq c_3 m. \quad (50)$$

Given $m > q + 2$ and a fixed choice x different from p, t , we note that the columns of column sum 3 cannot have the submatrix $q \cdot (\mathbf{1}_2 \mathbf{0}_0)$ in rows p, x nor the submatrix $q \cdot (\mathbf{1}_2 \mathbf{0}_0)$ in rows t, x since either would produce the configuration $q \cdot (\mathbf{1}_2 \mathbf{0}_2)$. Thus for a fixed $x \neq p, t$ (of which there are $m - 2$ choices), there are at most $2(q - 1)$ choices for j such that pair xj has triple overlapping pt in columns of column sum 3. A similar argument applies to the columns of column sum $m - 3$. Thus there exists a constant $c_4 = 2(q - 1)$ so that

$$|\mathcal{T}(pt)| \leq 2 \left(\frac{(m-2)2(q-1)}{2} \right) \leq c_4 m \quad (51)$$

Lemma 4.2 *There exists a constant $m_1 \geq q + 4$ so that for $m > m_1$, we have for all $\{i, j\} \subset [m]$ that $d_0(ij) + d_1(ij) \leq q - 3 + \epsilon(ij)$.*

Proof: Assume the contrary that pt is an index with $d_0(pt) + d_1(pt) \geq q - 3 + \epsilon(pt) + 1$. Let $\{r, s\} \subset [m] \setminus \{p, t\}$ and $rs \notin \mathcal{U}(pt) \cup \mathcal{T}(pt)$ (by (50) and (51) there are $\binom{m}{2} - c_3m - c_4m$ choices for rs). There are at most $2q - 2$ submatrices as in (49) contained in A else A has the configuration $q \cdot (\mathbf{1}_2\mathbf{0}_2)$. There are $4 - \epsilon(pt) - \epsilon(rs)$ submatrices (49) contained in columns of column sum 2, $m - 2$ and since $rs \notin \mathcal{U}(pt) \cup \mathcal{T}(pt)$ there are $(d_1(pt) + d_0(rs)) + (d_0(pt) + d_1(rs))$ submatrices (49) in columns of column sum 3, $m - 3$. Thus

$$(d_1(pt) + d_0(rs)) + (d_0(pt) + d_1(rs)) + 4 - \epsilon(pt) - \epsilon(rs) \leq 2(q - 1)$$

Substituting $d_0(pt) + d_1(pt) \geq q - 3 + \epsilon(pt) + 1$ and rearranging yields

$$d_0(rs) + d_1(rs) \leq (q - 3) - 1 + \epsilon(rs). \quad (52)$$

We wish to bound a_3 using (46). We split all pairs ij into three sets: those with $\{i, j\} \cap \{p, t\} = \emptyset$ and $ij \notin \mathcal{U}(pt) \cup \mathcal{T}(pt)$, those with $ij \in \mathcal{U}(pt) \cup \mathcal{T}(pt)$ (which forces $\{i, j\} \cap \{p, t\} = \emptyset$) and those with $\{i, j\} \cap \{p, t\} \neq \emptyset$. In the first case, we use (52).

$$\sum_{\substack{\{i,j\} \subset [m] \\ ij \notin \mathcal{U}(pt) \cup \mathcal{T}(pt) \\ \{i,j\} \cap \{p,t\} = \emptyset}} d_0(ij) + d_1(ij) \leq \sum_{\substack{\{i,j\} \subset [m] \\ ij \notin \mathcal{U}(pt) \cup \mathcal{T}(pt) \\ \{i,j\} \cap \{p,t\} = \emptyset}} (q - 3) - 1 + \epsilon(ij)$$

In the latter cases, note that $d_0(ij) \leq q - 1$ and $d_1(ij) \leq q - 1$ else, since $m \geq q + 4$, we would find a copy of $q \cdot (\mathbf{1}_2\mathbf{0}_2)$.

$$\begin{aligned} \sum_{\substack{\{i,j\} \subset [m] \\ ij \in \mathcal{U}(pt) \cup \mathcal{T}(pt)}}} d_0(ij) + d_1(ij) &\leq (c_3 + c_4)m \cdot 2(q - 1) \\ \sum_{\substack{\{i,j\} \subset [m] \\ \{i,j\} \cap \{p,t\} \neq \emptyset}} d_0(ij) + d_1(ij) &\leq 2(m - 2) \cdot 2(q - 1) \end{aligned}$$

Let c_5 be a constant chosen so that $c_5 > 2(c_3 + c_4 + 2)(q - 1)$. Combining yields

$$\sum_{ij} (d_0(ij) + d_1(ij)) \leq \sum_{\substack{\{i,j\} \subset [m] \\ ij \notin \mathcal{U}(pt) \cup \mathcal{T}(pt) \\ \{i,j\} \cap \{p,t\} = \emptyset}} \left((q - 3) - 1 + \epsilon(ij) \right) + c_5m.$$

Now using (40) and substituting for a_2 using (48) and substituting for a_3 using (46) and the above inequality with the estimate that there are at most $\binom{m-2}{2}$ choices for pairs ij with $\{i, j\} \cap \{p, t\} = \emptyset$ $ij \notin \mathcal{U}(ij) \cup \mathcal{T}(ij)$ and substituting for a_4 using (43):

$$2 \binom{m}{2} - (|E_0| + |E_1|) + \frac{(q - 3) - 1}{3} \binom{m - 2}{2} + \frac{1}{3} (|E_0| + |E_1|) + \frac{c_5}{3}m + c_2m > \binom{m}{2} \frac{q + 3}{3}$$

The coefficient of m^2 on the left side of the above inequality is only $\frac{q+2}{6}$ while on the right side is $\frac{q+3}{6}$. Thus there exists a constant m_1 so that for $m > m_1$, we have a contradiction proving the claim. \blacksquare

Let

$$Y = \{ij : d_0(ij) + d_1(ij) = q - 3 \text{ and } \epsilon(ij) = 0\}$$

Lemma 4.3 *There exists a constant c_6 so that*

$$|Y| > \binom{m}{2} - c_6 m. \quad (53)$$

Proof: We partition the $\binom{m}{2}$ pairs ij into 3 parts: Y , $E_0 \cup E_1$ and the rest. We note that for $ij \notin Y \cup E_0 \cup E_1$, we have $\epsilon(ij) = 0$ and $d_0(ij) + d_1(ij) \leq (q - 3) - 1$ by Lemma 4.2. Thus from (46) and using Lemma 4.2

$$\begin{aligned} a_3 = \frac{1}{3} \sum_{ij} d_0(ij) + d_1(ij) &\leq \frac{1}{3} ((q - 3)|Y| + ((q - 3) + 2)|E_0 \cup E_1| \\ &\quad + ((q - 3) - 1) \left(\binom{m}{2} - |Y| - |E_0 \cup E_1| \right)) \end{aligned}$$

Thus

$$a_3 \leq \frac{1}{3} \left((q - 3) \binom{m}{2} + 3|E_0 \cup E_1| - \binom{m}{2} + |Y| \right) \quad (54)$$

Using (48),(43), (54) in (40), we have

$$\begin{aligned} 2 \binom{m}{2} - (|E_0| + |E_1|) + \frac{1}{3} \left((q - 3) \binom{m}{2} + 3|E_0 \cup E_1| + \left(|Y| - \binom{m}{2} \right) \right) &+ c_2 m \\ &> \frac{q + 3}{3} \binom{m}{2}. \end{aligned}$$

We deduce, noting that $|E_0| + |E_1| \geq |E_0 \cup E_1|$, that $\frac{1}{3} (|Y| - \binom{m}{2}) + c_2 m > 0$ and so $|Y| > \binom{m}{2} - 3c_2 m$. Thus (53) holds for $c_6 = 3c_2$. \blacksquare

Form a graph G whose vertex set is the rows of the matrix A with edges ij for those $ij \in Y$. Thus G has at least $\frac{m^2}{2} - c_6 m$ edges. By Turán's Theorem [7], a graph with more than $\frac{m^2}{2} - \frac{m^2}{2(k-1)}$ edges has a clique of k vertices. Choose a constant c_7 so that for any choices $i, j \in [m]$

$$\begin{aligned} \left(\frac{c_7 \sqrt{m} - 2(q - 1)}{2} \right) &> (c_3 + c_4) m (> |\mathcal{T}(ij)| + |\mathcal{U}(ij)|), \quad \frac{1}{2} \left(\frac{c_7 \sqrt{m}}{2} \right) - 2m > (c_3 + c_4) m (> |\mathcal{T}(ij)| + |\mathcal{U}(ij)|), \\ \left(\frac{c_7 \sqrt{m}}{2} \right) &> (c_3 + c_4) m (\geq |\mathcal{T}(ij)| + |\mathcal{U}(ij)|). \end{aligned} \quad (55)$$

Then by Turán's Theorem, there exists a $M > m_0, m_1$ (m_0 is from Lemma 4.1 and m_1 is from Lemma 4.2) so that for $m > M$, graph G has a clique of $c_7\sqrt{m}$ vertices.

Let B denote the set of the rows in this clique. Hence for every $i, j \in B$ we have $d_1(ij) + d_0(ij) = q - 3$ and $\epsilon(ij) = 0$. Let A_4 denote the columns of A of column sum 4, 5, \dots , $m - 5$ or $m - 4$. Let $A_4(B)$ be the submatrix of A_4 of the rows indexed by B .

Lemma 4.4 *Assume $m > M$. Then $A_4(B)$ has no configuration $\mathbf{1}_2\mathbf{0}_2$.*

Proof: Assume there are rows $i, j, k, l \in B$ and a column α of A_4 with 0's in rows i, j and 1's in rows k, l . Without loss of generality, we may assume that there are more 1's than 0's in α in the rows of B so that the number of 1's in the rows of B is more than $c_7\sqrt{m}/2$. Thus by the third inequality in (55), we can find a pair gh of rows with $g, h \in B$, so that α has 1's in row g, h and $gh \notin \mathcal{T}(ij) \cup \mathcal{U}(ij)$. We may now argue that for our choice of i, j, g, h , we have $(d_1(ij) + d_0(gh)) + (d_1(gh) + d_0(ij)) + 4 - \epsilon(ij) - \epsilon(gh) = 2(q - 1)$ submatrices

$$\begin{array}{c} i \\ j \\ g \\ h \end{array} \begin{bmatrix} 1 \\ 1 \\ 0 \\ 0 \end{bmatrix} \quad \text{or} \quad \begin{array}{c} i \\ j \\ g \\ h \end{array} \begin{bmatrix} 0 \\ 0 \\ 1 \\ 1 \end{bmatrix} \quad (56)$$

in A in columns of column sum 2, 3, $m - 3$, $m - 2$. With another such submatrix in α in A_4 , we have $2(q - 1) + 1$ such submatrices, for our chosen quadruple i, j, g, h and so A has the configuration $q \cdot (\mathbf{1}_2\mathbf{0}_2)$, a contradiction. ■

Lemma 4.5 *Assume $m > M$. Then the inequality (9) holds.*

Proof: Assume $m > M$ and (40). Using Lemma 4.4, the columns of A_4 can be partitioned into two parts: Z the columns that have at most one 1 in the rows B and J the columns that have at most one 0 in the rows of section B .

For each pair $i, j \in [m] \setminus B$, let $t(ij)$ count the sum of the number of columns in Z with 1's in both rows i, j as well as the number of columns in J with 0's in both rows i, j . For all other pairs ij , let $t(ij) = 0$. Given the column sums in A_4 , every column in Z has at least three 1's in rows $[m] \setminus B$ and every column in J has at least three 0's in rows $[m] \setminus B$. We have

$$\sum_{ij} t(ij) \geq 3a_4 \quad (57)$$

Moreover, we find that $t(ij) \leq 2(q - 1)$: Given a choice for i, j , if we have q columns in Z with 1's in rows i, j then there are at most q rows of B containing 1's for these q columns (since each column of Z has at most one 1 in the rows of B). But then if we choose two rows of B from the remaining $\geq |B| - q$ rows in conjunction with i, j then we have a copy of the configuration $q \cdot (\mathbf{1}_2\mathbf{0}_2)$. Similarly, there cannot be q columns of J with 0's on rows i, j . We conclude $t(ij) \leq 2(q - 1)$.

For a given pair $i, j \in [m] \setminus B$, consider the $t(ij)$ columns contributing to $t(ij)$. By the first inequality in (55), we can find a pair of rows gh ($g, h \in B$) so that $gh \notin \mathcal{T}(ij) \cup \mathcal{U}(ij)$

and in addition g, h are not chosen from the up to $2(q-1)$ rows of B which are given as follows: the $\leq q-1$ rows of B which have 1's in the columns of Z having 1's in both rows i, j and the $\leq q-1$ rows of B which have 0's in the columns of J having 0's in both rows i, j . Thus if α is a column of Z with 1's in rows i, j then α has 0's in rows g, h and if α is a column of J with 0's in rows i, j then α has 1's in rows g, h . There will be $4 - \epsilon(ij) - \epsilon(gh)$ submatrices as in (56) in the columns of column sum 2 or $m-2$. Neither pair ij has triple overlapping gh nor pair gh has triple overlapping ij and so there will be $(d_1(ij) + d_0(gh)) + (d_0(ij) + d_1(gh))$ submatrices as in (56) in the columns of column sum 3 or $m-3$. By our choice of g, h , a column α in Z with 1's in rows i, j will have 0's on rows g, h . A column β in J with 0's in rows i, j will have 1's on rows g, h . Thus in A_4 we can find $t(ij)$ submatrices as in (56). In the matrix A , an ordered quadruple of rows i, j, g, h has at most $2(q-1)$ submatrices as given in (56) else A would have the configuration $q \cdot (\mathbf{1}_2 \mathbf{0}_2)$. Thus

$$(d_1(ij) + d_0(gh)) + (d_0(ij) + d_1(gh)) + 4 - \epsilon(ij) - \epsilon(gh) + t(ij) \leq 2(q-1).$$

Substituting $d_0(gh) + d_1(gh) = q-3$ and $\epsilon(gh) = 0$ and rearranging we have

$$d_1(ij) + d_0(ij) \leq (q-3) + \epsilon(ij) - t(ij). \quad (58)$$

This inequality is true for other i, j using Lemma 4.2 when $t(ij) = 0$. Thus

$$\sum_{ij} (d_0(ij) + d_1(ij)) \leq \sum_{ij} (q-3 + \epsilon(ij) - t(ij)) \quad (59)$$

Taking (40) with a_2 from (48) and with a_3 from (46) using (58) we obtain

$$2 \binom{m}{2} - |E_0| - |E_1| + \frac{1}{3} \sum_{ij} (q-3 + \epsilon(ij) - t(ij)) + a_4 > \frac{q+3}{3} \binom{m}{2}$$

Simplifying and using $\sum_{ij} \epsilon(ij) = |E_0| + |E_1|$ and (57) we obtain

$$-\frac{2}{3}(|E_0| + |E_1|) > 0$$

which is a contradiction (even for $|E_0| + |E_1| = 0$). This establishes (9). \blacksquare

Proof of Proposition 1.8: Lemma 4.5 establishes most of Proposition 1.8 but we are also interested in cases when the bound is achieved. Assume $m > M$ and $m \equiv 1, 3 \pmod{6}$. We now consider an m -rowed simple matrix A which has no configuration $q \cdot (\mathbf{1}_2 \mathbf{0}_2)$ and with $\binom{m}{0} + \binom{m}{1} + \frac{q+3}{3} \binom{m}{2} + \binom{m}{m-1} + \binom{m}{m}$ columns. One repeats the previous lemmas and arguments replacing the inequality (40) with the equation

$$a_2 + a_3 + a_4 = \frac{q+3}{3} \binom{m}{2}. \quad (60)$$

We wish to show $a_2 = 2\binom{m}{2}$, $a_4 = 0$, $a_3 = \frac{q-3}{3}\binom{m}{2}$ and there exists positive integers a, b , $a+b = q-3$ so that for all pairs ij , $d_0(ij) = a$ and $d_1(ij) = b$. Now Lemma 4.1 holds with (40) as an equality. We deduce the same bounds for $\mathcal{U}(pt)$ and $\mathcal{T}(pt)$. Lemma 4.2 still holds since the final contradiction does not require the strict inequality of (40) merely the equality of (60). Lemma 4.3 holds and we can choose B as large as possible but at least satisfying the inequalities (55). Lemma 4.4 continues to hold.

Assume that not all pairs pt with $p, t \in B$ have the same value for $d_0(pt)$. We can choose ij with $i, j \in B$ so that at least $\frac{1}{2}\binom{|B|}{2}$ pairs pt of $\binom{B}{2}$ have $d_0(ij) \neq d_0(pt)$. Then the number of pairs pt of $\binom{B}{2}$ in $\binom{B \setminus \{i, j\}}{2}$ with $d_0(ij) \neq d_0(pt)$ is at least $\frac{1}{2}\binom{|B|}{2} - 2|B|$. Now using the second inequality of (55) with $|\mathcal{U}(ij)| + |\mathcal{T}(ij)| \leq (c_3 + c_4)m$ and $|B| \leq m$, we can find a pair $k, l \in B \setminus \{i, j\}$ with $d_0(ij) \neq d_0(kl)$, $kl \notin \mathcal{U}(ij) \cup \mathcal{T}(ij)$. By definition of B ,

$$d_0(ij) + d_1(ij) = q - 3, \quad d_0(kl) + d_1(kl) = q - 3.$$

We may assume without loss of generality that $d_0(kl) < d_0(ij)$, $d_1(kl) > d_1(ij)$ and then

$$d_0(ij) + d_1(kl) \geq q - 2$$

We also have $\epsilon(ij) = \epsilon(kl) = 0$. Then A has a column of column sum 2 and a column of column sum $m - 2$ both with 1's in rows k, l and 0's in rows i, j . Also we have $d_0(ij) + d_1(kl)$ columns with 1's in rows k, l and 0's in rows i, j since $kl \notin \mathcal{U}(ij) \cup \mathcal{T}(ij)$. But then A has $q \cdot (\mathbf{1}_2 \mathbf{0}_2)$, a contradiction. We conclude that all pairs pt with $p, t \in B$ have the same value for $d_0(pt)$.

We follow our proof of Lemma 4.5 using (60) and deduce that $E_0 \cup E_1 = \emptyset$ and so $a_2 = 2\binom{m}{2}$. Also we deduce that

$$\sum_{ij} t(ij) = 3a_4$$

and as a result we can deduce that any column α in A_4 either has column sum 4 with exactly one 1 in a row of B or has column sum $m - 4$ with exactly one 0 in a row of B .

Assume $a_4 > 0$ and consider α in A_4 , say with column sum 4 and with 1's in rows i, j, k, l where $i \in B$ and $j, k, l \in \{1, 2, \dots, m\} \setminus B$. Choose $r, s \in B \setminus i$ so that $d_0(rs) + d_1(rs) = q - 3$ and with $rs \notin \mathcal{T}(ij) \cup \mathcal{U}(ij)$ (using first inequality of (55)). Column α has 1's in rows i, j and 0's in row r, s . Using $E_0 \cup E_1 = \emptyset$, we deduce that $d_1(ij) + d_0(rs) \leq q - 3 - 1$ and $d_0(ij) + d_1(rs) \leq q - 3$ else if either inequality is violated we create $q \cdot (\mathbf{1}_2 \mathbf{0}_2)$. We deduce $d_0(ij) + d_1(ij) \leq (q - 3) - 1$. This yields a slight variant of (59):

$$\sum_{ij} (d_0(ij) + d_1(ij)) \leq \sum_{ij} (q - 3 + \epsilon(ij) - t(ij)) - 1.$$

The extra '-1' is sufficient to obtain a contradiction when we substitute for a_2, a_3, a_4 in (60). We then deduce $a_4 = 0$.

With $a_4 = 0$ and $a_2 = 2\binom{m}{2}$, we deduce $a_3 = \frac{q-3}{3}\binom{m}{2}$ using (60). Given that $\epsilon(ij) = 0$ for all ij and using Lemma 4.2, we deduce $d_0(ij) + d_1(ij) = q - 3$ for all pairs ij and so $B = \{1, 2, \dots, m\}$. Our above arguments tell us $d_0(pt)$ is the same for every choice $p, t \in B$, allowing us to conclude that there exists positive integers a, b , $a + b = q - 3$ so that for all pairs ij , $d_0(ij) = a$ and $d_1(ij) = b$. From this we can readily conclude that the columns of column sum 3 correspond to a 2-design $S_a(2, 3, m)$ and the columns of column sum $m - 3$ correspond to the (0,1)-complement of a 2-design $S_b(2, 3, m)$. ■

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